

# **ECE 2400 Computer Systems Programming Fall 2017**

## **Topic 18: Tables**

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revision: 2017-11-29-12-14

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## 1. Tables Concepts

- Tables have rows and columns
- Rows are usually identified by a name (number or even a string)
- Columns are usually identified by a name (number or even a string)
- Rows may have multiple columns of data
- Tables are indexed to access row data at a certain column

|   | A | B | C |
|---|---|---|---|
| 0 |   |   |   |
| 1 |   |   |   |
| 2 |   |   |   |

Table as a Set ADT

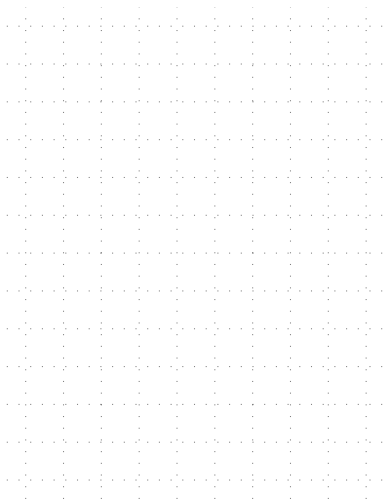


Table as a Map ADT



## 2. Lookup Tables

- A lookup table is a table in which all potential data is known
- A lookup table is accessed by simple array indexing in  $O(1)$  time
- Lookup tables may be multidimensional

### 2.1. Set of Integers as a Lookup Table

- A set of unsigned integers in the range  $[0, 4]$
- Implemented internally as a simple array of booleans

| Index    | 0 | 1 | 2 | 3 | 4 |
|----------|---|---|---|---|---|
| Present? |   |   |   |   |   |

#### Interface

```
1  class LookupTableInt
2  {
3  public:
4      LookupTableInt();
5
6      void add      ( unsigned int v );
7      void remove   ( unsigned int v );
8      bool contains ( unsigned int v );
9
10     private:
11         static const unsigned int size = 5;
12         bool m_table[size];
13     };
```

## Implementation

```
1 LookupTableInt::LookupTableInt()
2 { }
3
4 void LookupTableInt::add( unsigned int v )
5 {
6     m_table[v] = true;
7 }
8
9 void LookupTableInt::remove( unsigned int v )
10 {
11     m_table[v] = false;
12 }
13
14 bool LookupTableInt::contains( unsigned int v )
15 {
16     return m_table[v];
17 }
```

## Main

```
1 int main( void )
2 {
3     LookupTableInt table;
4
5     table.add(3);    // Add value 3
6     table.remove(3); // Remove value 3
7
8     return 0;
9 }
```

## 2.2. Set of Strings as a Lookup Table

- A set of strings limited to "Apple", "Banana", "Cherry"
- Implemented internally as a simple array of booleans
- Must correspond each possible string to an internal array index

| Index    | Apple | Banana | Cherry |
|----------|-------|--------|--------|
| Present? |       |        |        |

### Interface

```
1  class LookupTableString
2  {
3      public:
4          LookupTableString();
5
6          void add      ( std::string v );
7          void remove   ( std::string v );
8          bool contains ( std::string v );
9
10         private:
11             unsigned int ind( std::string v );
12             static const unsigned int size = 3;
13             bool m_table[size];
14     };
```

### Implementation

```
1  LookupTableString::LookupTableString()
2  { }
```

```
4 void LookupTableString::add( std::string v )
5 {
6     m_table[ ind(v) ] = true;
7 }
8
9 void LookupTableString::remove( std::string v )
10 {
11     m_table[ ind(v) ] = false;
12 }
13
14 bool LookupTableString::contains( std::string v )
15 {
16     return m_table[ ind(v) ];
17 }
18
19 unsigned int LookupTableString::ind( std::string v )
20 {
21     if ( v == "Apple" ) return 0;
22     else if ( v == "Banana" ) return 1;
23     else if ( v == "Cherry" ) return 2;
24     else throw std::out_of_range( "No such fruit!" );
25 }
```

## Main

```
1 int main( void )
2 {
3     LookupTableString table;
4
5     table.add( "Cherry" ); // Add value "Cherry"
6     table.remove( "Cherry" ); // Remove value "Cherry"
7
8     return 0;
9 }
```

## 2.3. Strengths and Weaknesses of Lookup Tables

- We have implemented a set of integers as a lookup table
- We have implemented a set of strings as a lookup table
- Recall that lists, vectors, and even trees can also implement sets!

Set ADT as a List



Set ADT as a Tree



|             | add          | remove       | contains     |
|-------------|--------------|--------------|--------------|
| list/vector | $O(N)$       | $O(N)$       | $O(N)$       |
| tree        | $O(\log(N))$ | $O(\log(N))$ | $O(\log(N))$ |
| lut         | $O(1)$       | $O(1)$       | $O(1)$       |

### Strengths of Lookup Tables

- Lookup tables have great time complexity!

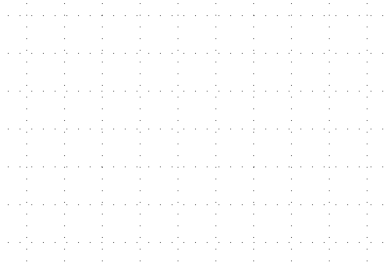
### Weaknesses of Lookup Tables

- I want a set for a larger range of numbers (e.g., from 0 to 1,000,000)
  - Lookup tables are \_\_\_\_\_
- I want a set of other types (e.g., strings)
  - Lookup tables are \_\_\_\_\_

## 2.4. Mitigating the Weaknesses of Lookup Tables

- Lookup tables have space complexity that is linear with  $k$ , the space of all possible unique inputs.

$k = 4$  and  $N = 2$



$k = 1,000,000$  and  $N = 2$



- How can we make the space complexity instead scale with  $N$ , the number of unique elements currently in the set?

### Reduction Operations

- Define a `reduce` function that converts a large number into the range accessible by the internal array.
- When the internal array fills up, double the size of the array.

### Member function to reduce a large number

```
1 unsigned int reduce( unsigned int v ) {
```

---

---

---

---

---

## Interface

```
1  class LookupTableIntOneMillion
2  {
3      public:
4          LookupTableIntOneMillion();
5
6          void add      ( unsigned int v );
7          void remove   ( unsigned int v );
8          bool contains ( unsigned int v );
9
10     private:
11         unsigned int reduce ( unsigned int v );
12         static const unsigned int size = 10;
13         bool m_table[size];
14 };
```

## Implementation

```
1  LookupTableIntOneMillion::LookupTableIntOneMillion()
2  { }
3
4  void LookupTableIntOneMillion::add( unsigned int v )
5  {
6      m_table[ reduce(v) ] = true;
7  }
8
9  void LookupTableIntOneMillion::remove( unsigned int v )
10 {
11     m_table[ reduce(v) ] = false;
12 }
13
14 bool LookupTableIntOneMillion::contains( unsigned int v )
15 {
16     return m_table[ reduce(v) ];
17 }
```

## Implementation of New Member Function

```
18 unsigned int LookupTableIntOneMillion::reduce( unsigned int v
19 {
20     return v % size;
21 }
```

## Main

```
1 int main( void )
2 {
3     LookupTableIntOneMillion table;
4
5     table.add( 967 ); // Reduces to index 7
6     table.add( 333 ); // Reduces to index 3
7     table.add( 555 ); // Reduces to index 5
8
9     return 0;
10 }
```

## Space Complexity with a Reduction Function

- Ignoring conflicts, we now have space complexity linear with N
- Time complexity remains at  $O(1)$  access
- Achieving this space complexity cost us...
  - Increased work in order to double the size of the internal array
  - Increasing the work on each access
- Any other issues?

## 2.5. Mitigating the Weaknesses of Lookup Tables (cont.)

- Lookup tables are *inflexible* in storing non-integer types like strings.
- But strings, like any other type represented by a computer, are actually just numbers in the underlying representation!
- The key idea is to add a new function, a *hash* function, which converts a string or other arbitrary type into an integer.
  - Values of `char` type are just numbers!
  - The ASCII char 'a' is 97, 'b' is 98, etc...

### Member function to convert a char array into an unsigned integer

```
1 unsigned int hash( char v[] ) {
```

---

---

---

---

---

---

## Good Hash Functions

- What makes a hash function a “good” hash function?
- Property 1: We want a *valid* hash function
  - Returns the same value on subsequent calls to the same item
  - For any equivalent objects  $a == b$ , their hashes are also equal



- Property 2: We want a hash function that provides *uniformity*
  - Maps the expected inputs as evenly as possible over the output range
  - Specifically, the hash result should not be a value (e.g., 100) more often



- Property 3: We want a hash function with  $O(1)$  time complexity

## Interface

```
1  class LookupTableStringHash
2  {
3      public:
4          LookupTableStringHash();
5          void add      ( std::string v );
6          void remove   ( std::string v );
7          bool contains ( std::string v );
8
9      private:
10         unsigned int reduce ( unsigned int v );
11         unsigned int hash   ( std::string v );
12         static const unsigned int size = 10;
13         bool m_table[size];
14     };
```

## Implementation

```
1  LookupTableStringHash::LookupTableStringHash()
2  { }
3
4  void LookupTableStringHash::add( std::string v )
5  {
6      m_table[ reduce( hash(v) ) ] = true;
7  }
8
9  void LookupTableStringHash::remove( std::string v )
10 {
11     m_table[ reduce( hash(v) ) ] = false;
12 }
13
14 bool LookupTableStringHash::contains( std::string v )
15 {
16     return m_table[ reduce( hash(v) ) ];
17 }
```

## Implementation of New Member Functions

```
1 unsigned int LookupTableStringHash::reduce( unsigned int v )
2 {
3     return v % size;
4 }
5
6 unsigned int LookupTableStringHash::hash( std::string v )
7 {
8     unsigned int hashnum = 0;
9     for ( auto& c : v )
10         hashnum += c;
11     return hashnum;
12 }
```

## Main

```
1 int main( void )
2 {
3     LookupTableStringHash table;
4
5     table.add( "c" );    // Hash 99 : Reduced to 9
6     table.add( "a" );    // Hash 97 : Reduced to 7
7     table.add( "t" );    // Hash 116 : Reduced to 6
8
9     table.add( "cat" );  // Hash: 99 + 97 + 116 = 312
10                        // Reduced to 2
11
12     return 0;
13 }
```

## Summary

- We started off trying to mitigate the weaknesses of lookup tables

|                      | add          | remove       | contains     |
|----------------------|--------------|--------------|--------------|
| list/vector          | $O(N)$       | $O(N)$       | $O(N)$       |
| tree                 | $O(\log(N))$ | $O(\log(N))$ | $O(\log(N))$ |
| lut                  | $O(1)$       | $O(1)$       | $O(1)$       |
| lut w/ hash & reduce | $O(1)$       | $O(1)$       | $O(1)$       |

## Strengths of Lookup Tables with hash and reduce Functions

- Time complexity is still  $O(1)$

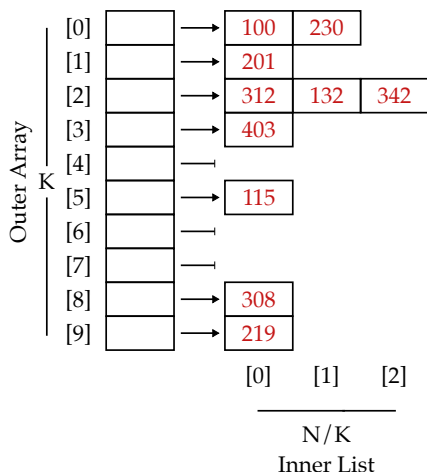
## Mitigated Weaknesses of Lookup Tables with hash and reduce

- I want a set for a larger range of numbers (e.g., from 0 to 1,000,000)
  - Lookup tables were *space-inefficient*
  - Lookup tables with hash and reduce have space efficiency of  $O(N)$
- I want a set of other types (e.g., strings)
  - Lookup tables were *inflexible*
  - Lookup tables with hash and reduce can handle any type as long as it provides a valid hash function
- What do we do about collisions?

### 3. Hash Tables

- Hash tables are lookup tables with hash and reduce functions that also handle collisions.
- Two common approaches for handling collisions
  - Separate chaining (usually with linked lists)
  - Open addressing (usually with linear probing)

#### 3.1. Collision Handling with Separate Chaining



- The internal array elements are lists in a separate chaining scheme
- With a good hash function, the  $N$  elements are spread evenly over the  $K$  buckets.
- Adding a new element to the hash table is now a push back on a list
- What does this do to our time complexity?

```
1  class List
2  {
3      public:
4
5          class Itr
6          {
7              public:
8                  Itr( Node* node_p );
9                  void      next();
10                 std::string& get();
11                 bool      eq( const Itr& itr ) const;
12
13             private:
14                 friend class List;
15                 Node* m_node_p;
16         };
17
18         Itr begin();
19         Itr end();
20
21         List();
22         ~List();
23         void push_front( std::string& v );
24
25     private:
26
27         struct Node
28         {
29             std::string value;
30             Node*      next_p;
31         };
32
33         Node* m_head_p;
34     };
```

```
1 List::Itr    operator++( List::Itr& itr, std::string );
2 List::Itr&   operator++( List::Itr& itr );
3 std::string& operator* ( List::Itr& itr );
4 bool        operator!=( const List::Itr& itr0,
5                          const List::Itr& itr1 );
```

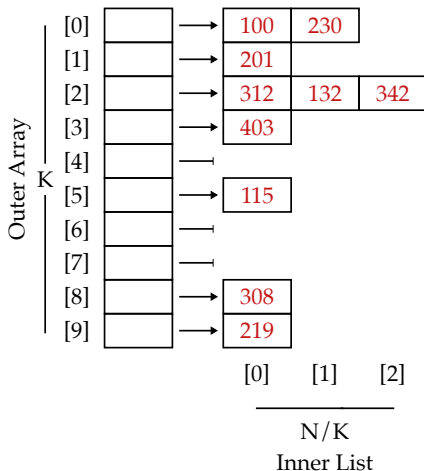
- With the above List class (provided exactly as discussed in previous lectures), implement the add function of a hash table that handles collisions with chaining using linked lists.
- Note that a templated argument Hasher is often used to provide a class with a hash function.
  - The following class “has-a” Hasher named hash.
  - The Hasher has an overloaded call operator (i.e., operator()) that does the hash (e.g., calling hash with a string returns an unsigned integer).
- Below is the interface of the hash table. Assume that calling the Hasher object using hash() will create some collisions.

```
1 template < typename Hasher >
2 class HashTableChaining
3 {
4     public:
5         HashTableChaining();
6         void add      ( const std::string& v );
7         void remove   ( const std::string& v );
8         bool contains ( const std::string& v );
9
10        private:
11            unsigned int reduce( unsigned int v );
12            Hasher hash;
13            static const unsigned int size = 10;
14            List m_table[size];
15 };
```

[illegible]

## Summary

- How do hash tables compare to previous implementations of sets?



|                      | add          | remove       | contains     |
|----------------------|--------------|--------------|--------------|
| list/vector          | $O(N)$       | $O(N)$       | $O(N)$       |
| tree                 | $O(\log(N))$ | $O(\log(N))$ | $O(\log(N))$ |
| lut                  | $O(1)$       | $O(1)$       | $O(1)$       |
| lut w/ hash & reduce | $O(1)$       | $O(1)$       | $O(1)$       |
| hash table           | _____        | _____        | _____        |