

# **ECE 2400 Computer Systems Programming Fall 2017**

## **Topic 8: Algorithm Analysis**


School of Electrical and Computer Engineering  
Cornell University

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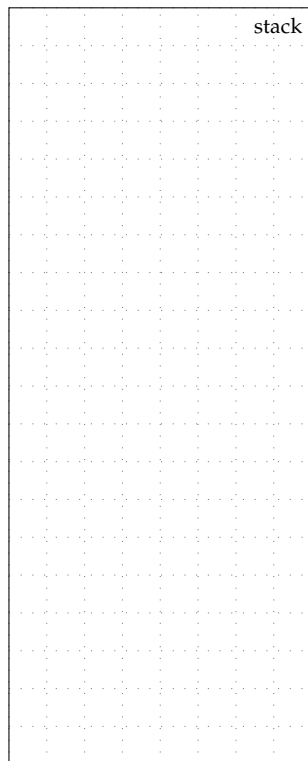
## 1. Analyzing Two Search Algorithms

- Assume we have a sorted input array of integers
- Consider algorithms to find if a given value is in the array
- The algorithm should return 1 if value is in array, otherwise return 0

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- Let  $N$  be the size of the input array
  - Let  $T$  be the execution time of an algorithm
  - Let  $S$  be the additional storage (space) required by the algorithm
  - Our goal is to derive equations for  $T$  and  $S$  as a function of  $N$
  - Our equations can be rough estimates
  - Execution time can be measured in number of C statements
  - Space requirements can be measured in number of C variables

## 1.1. Linear Search

```
1  int find( int a[], size_t size, int v )
2  {
3      for ( size_t i = 0; i < size; i++ ) {
4          if ( a[i] == v )
5              return 1;
6          // else if ( a[i] > v )
7          //     return 0;
8      }
9      return 0;
10 }
11
12 int main( void )
13 {
14     int a[] = { 0, 2, 4, 6, 8, 10, 12, 14 };
15     int find4  = find( a, 8,  4 );
16     int find0  = find( a, 8,  0 );
17     int find20 = find( a, 8, 20 );
18     return 0;
19 }
```

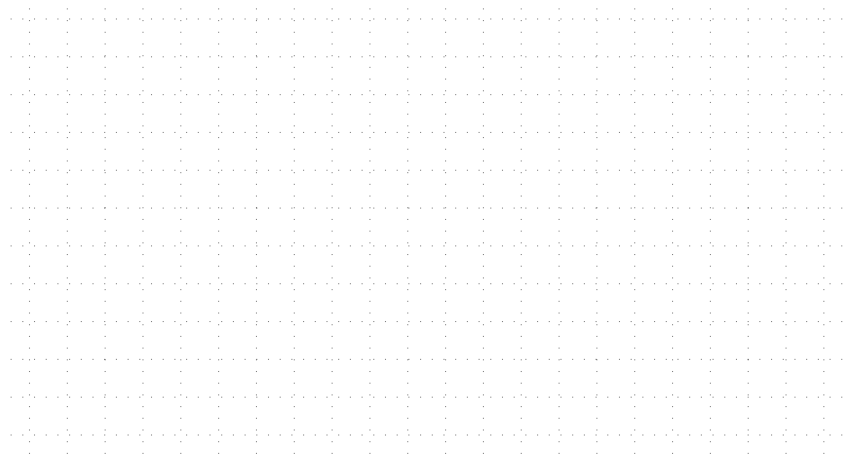


## 1.2. Binary Search

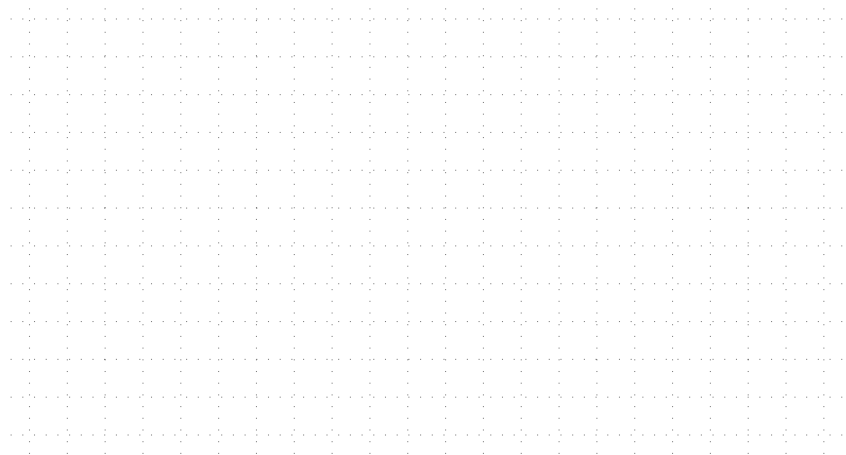
```
1  int find_h( int a[], size_t left,
2             size_t right, int v )
3  {
4      if ( a[left] == v )      return 1;
5      if ( a[right] == v )    return 1;
6      if ( (right-left) == 1 ) return 0;
7
8      int middle = left + (right-left)/2;
9      if ( a[middle] > v )
10         return find_h( a, left, middle, v );
11     else
12         return find_h( a, middle, right, v );
13 }
14
15 int find( int a[], size_t size, int v )
16 {
17     return find_h( a, 0, size-1, v );
18 }
19
20 int main( void )
21 {
22     int a[] = { 0, 2, 4, 6, 8, 10, 12, 14 };
23     int find4  = find( a, 8, 4 );
24     int find0  = find( a, 8, 0 );
25     int find20 = find( a, 8, 20 );
26     return 0;
27 }
```

stack

## Annotating call tree with execution time



## Annotating call tree with space requirements



### 1.3. Comparing Linear vs. Binary Search

Linear Search

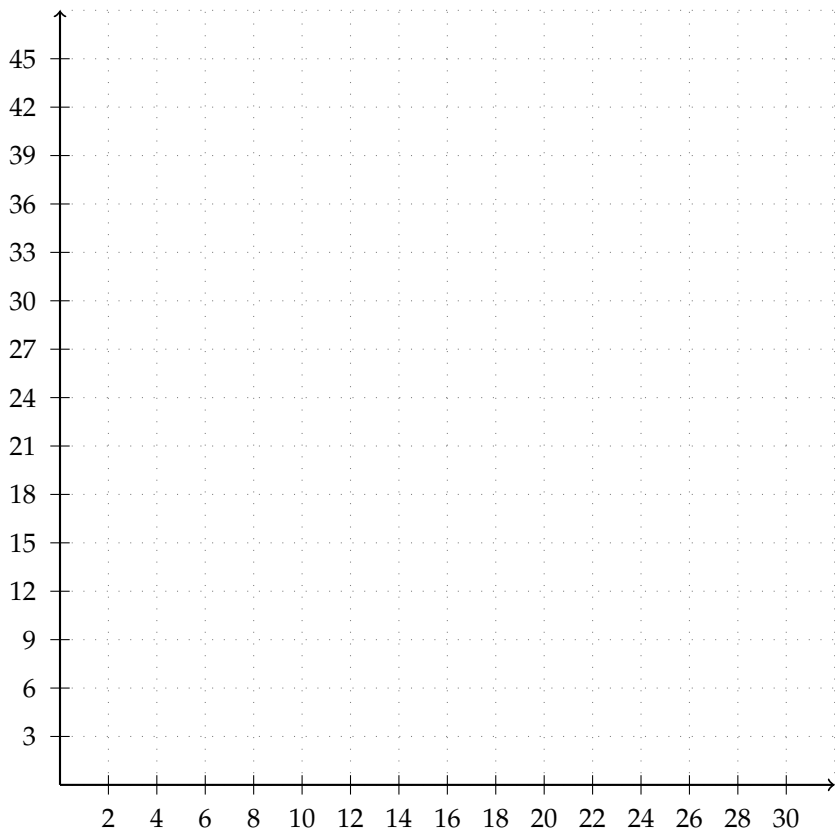
$$T_b(N) = 3$$

$$T_w(N) = 2N + 2$$

Binary Search

$$T_b(N) = 3$$

$$T_w(N) = 8 \log_2(N) + 3$$



Linear Search

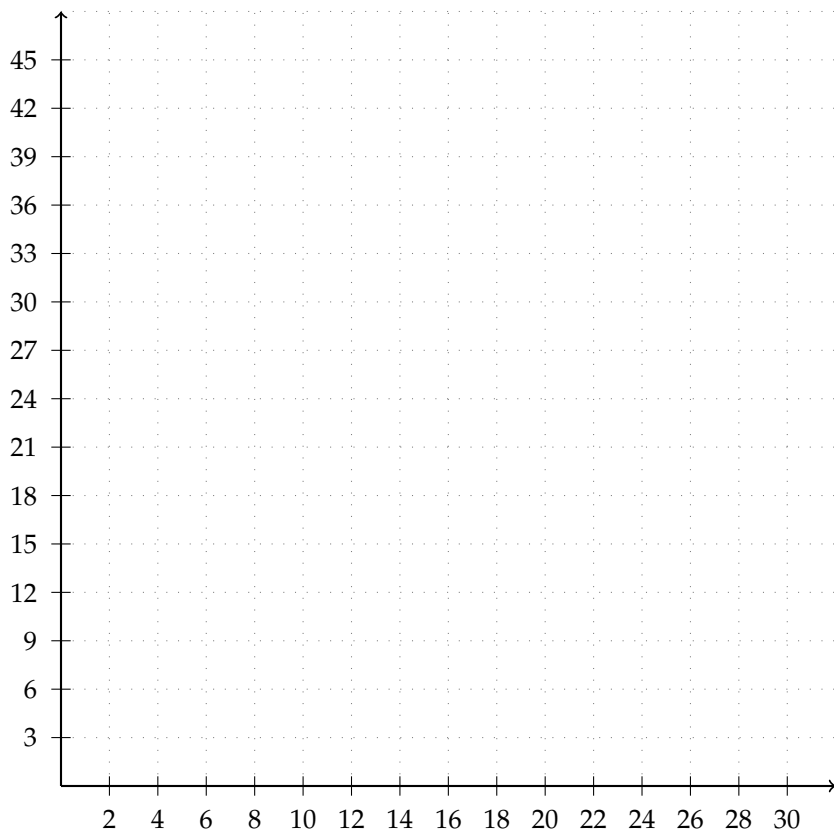
$$S_b(N) = 4$$

$$S_w(N) = 4$$

Binary Search

$$S_b(N) = 8$$

$$S_w(N) = 5 \log_2(N) + 4$$



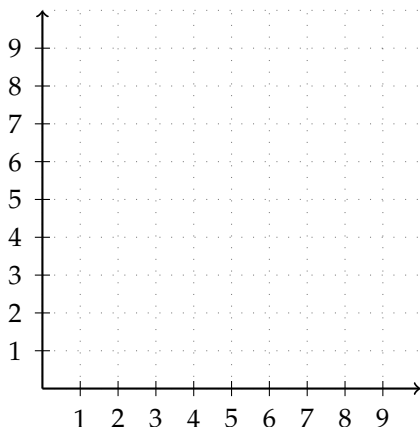


## 2. Time and Space Complexity

- We want a way to characterize algorithms at a high-level so we can quickly compare and contrast the expected performance of two algorithms as  $N$  grows large
- We want to gloss over low-level details
  - Absolute time of each C statement
  - Absolute storage requirements for each C variable
- Big-Oh notation captures the asymptotic behavior of a function

$$f(x) \text{ is } O(g(x)) \Leftrightarrow \exists x_0, c. \forall x > x_0. f(x) \leq c \cdot g(x)$$

- Formally:  $f(x)$  is  $O(g(x))$  if there is some value  $x_0$  and some value  $c$  such that for all  $x$  greater than  $x_0$ ,  $f(x) \leq c \cdot g(x)$
- Informally:  $f(x)$  is  $O(g(x))$  if  $g(x)$  captures the “most significant trend” of  $f(x)$  as  $x$  becomes very large



**Big-Oh Examples**

$f(x)$	is	$O(g(x))$
3	is	$O(1)$
$2N$	is	$O(N)$
$2N + 3$	is	$O(N)$
$4N^2$	is	$O(N^2)$
$4N^2 + 2N + 3$	is	$O(N^2)$
$4 \log_2(N)$	is	$O(\log_2(N))$
$N + 4 \log_2(N)$	is	$O(N)$

- Constant factors do not matter in big-oh notation
- Non-leading terms do not matter in big-oh notation
- What matters is the general trend as  $N$  becomes very large

**Big-Oh Classes**

	Class	N = 100 requires
$O(1)$	Constant Time	1 step
$O(\log_2(N))$	Logarithmic Time	6–7 steps
$O(N)$	Linear Time	100 steps
$O(N \cdot \log_2(N))$	Linearithmic Time	664 steps
$O(N^2)$	Quadratic Time	10K steps
$O(N^3)$	Cubic Time	1M steps
$O(N^c)$	Polynomial Time	
$O(2^N)$	Exponential Time	1e30 steps
$O(N!)$	Factorial Time	9e157 steps

- Exponential and factorial time algorithms are considered intractable
- With one nanosecond steps, exponential time would require many centuries and factorial time would require the lifetime of the universe

### Revisiting linear vs. binary search

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Linear	$T_w(N) = 2N + 2$	is $O(N)$	linear time
Binary	$T_w(N) = 8 \log_2(N) + 3$	is $O(\log_2(N))$	logarithmic time
<hr/>			
Linear	$S_w(N) = 4$	is $O(1)$	constant space
Binary	$S_w(N) = 5 \log_2(N) + 4$	is $O(\log_2(N))$	logarithmic space

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- Does this mean binary search is always faster?
- Does this mean linear search always require less storage?
- For very large  $N$ , but we don't always know  $x_0$ 
  - $T$  can have very large constants
  - $T$  can have non-leading terms
- This analysis is for worst case complexity
  - results can look very different for best case complexity (both  $O(1)$ )
  - results can look very different for typical/average complexity
- For reasonable problem sizes and/or different input data characteristics, sometimes an algorithm with worse time (space) complexity can still be faster (smaller)

### 3. Analysis of List and Vector Data Structures

Let's analyze the time and space complexity of various operations on the list and vector data structures.

#### Analysis of time and space complexity of `list_insert`

```
1 void list_push_front( list_t* list_p, int v )
2     allocate new node
3     set new node's value to v
4     set new node's next ptr to head ptr
5     set head ptr to point to new node
6
7 void list_insert( list_t* list_p, node_t* node_p, int v )
8     if list is empty
9         list_push_front( list_p, v )
10    else
11        allocate new node
12        set new node's value to v
13        set new node's next ptr to node's next ptr
14        set node's next ptr to point to new node
```

What is the time complexity for `list_insert`?

What is the space complexity for `list_insert`?

#### Analysis of time and space complexity of vector\_insert

```
1 void vector_push_front( vector_t* vec_p, int v )
2     set prev value to v
3     for i in 0 to vector's size (inclusive)
4         set temp value to vector's data[i]
5         set vector's data[i] to prev value
6         set prev value to temp value
7     set vector's size to size + 1
8
9 void vector_insert( vector_t* vec_p, size_t idx, int v )
10    if vector is empty
11        vector_push_front( vec_p, v )
12    else
13        set prev value to v
14        for i in idx+1 to vector's size (inclusive)
15            set temp value to vector's data[i]
16            set vector's data[i] to prev value
17            set prev value to temp value
18        set vector's size to size + 1
```

What is the time complexity for vector\_insert?

What is the space complexity for vector\_insert?

#### **Analysis of time and space complexity of list\_sorted\_insert**

```
1 void list_sorted_insert( list_t* list_p, int v )
2     set prev node ptr to head ptr
3     set curr node ptr to head node's next ptr
4
5     while curr node ptr is not NULL
6         if v is less than curr node's value
7             list_insert( list_p, prev node ptr, v )
8             return
9
10    set prev node ptr to curr node ptr
11    set curr node ptr to curr node's next ptr
```

What is the time complexity for list\_sorted\_insert?

What is the space complexity for list\_sorted\_insert?

#### **Analysis of time and space complexity of vector\_sorted\_insert**

```
1 void vector_sorted_insert( vector_t* vec_p, int v )
2     for i in 0 to vector's size
3         if v is less than vector's data[i]
4             vector_insert( vec_p, i-1, v )
5             return
```

What is the time complexity for vector\_sorted\_insert?

What is the space complexity for vector\_sorted\_insert?

#### **Analysis of time and space complexity of list\_sort\_insert**

```
1 void list_sort( list_t* list_p )
2   construct output list
3
4   set curr node ptr to input list's head ptr
5   while curr node ptr is not NULL
6     list_sorted_insert( output list, curr node's value )
7     set curr node ptr to curr node's next ptr
8
9   destruct input list
10  set input list's head ptr to output list's head ptr
```

What is the time complexity for list\_sort?

What is the space complexity for list\_sort?

#### **Analysis of time and space complexity of vector\_sort\_insert**

```
1 void vector_sort( vector_t* vec_p )
2   construct output vector
3
4   for i in 0 to vector's size
5     vector_sorted_insert( output vector, input vector's data[i] )
6
7   destruct input vector
8   set input vectors data ptr to output list's data ptr
```

What is the time complexity for vector\_sort?

What is the space complexity for vector\_sort?

Operation	Time Complexity		Space Complexity	
	List	Vector	List	Vector
insert				
sorted_insert				
sort				
push_front				
push_back				
find				

- Does this mean `list_sort` and `vector_sort` will have the same execution time? **absolutely not!**
- If two algorithms have the same time complexity, the constants and other terms are what makes the difference!
- This analysis is for worst case complexity
  - results can look very different for best case complexity
  - results can look very different for typical/average complexity
- In computer systems programming, we care about time and space complexity, but we also care about absolute execution time and absolute space requirements on a variety of inputs