

ECE 2300
Digital Logic & Computer Organization
Spring 2025

Combinational Logic Minimization



Cornell University

Announcements

- **Weekly calendar (including TA OH schedule) is posted on the course web**

De Morgan's Theorem

- Very important, also known as De Morgan's *Law*

$$(T12) \quad (X1 \cdot X2 \cdot \dots \cdot Xn)' = X1' + X2' + \dots + Xn'$$

$$(T12') \quad (X1 + X2 + \dots + Xn)' = X1' \cdot X2' \cdot \dots \cdot Xn'$$

De Morgan Example

- By DeMorgan's Law

$$(X \cdot Y \cdot Z)' = X' + Y' + Z'$$

- Proof by perfect induction

XYZ	$(X \cdot Y \cdot Z)'$	$X' + Y' + Z'$
000	1	1
001	1	1
010	1	1
011	1	1
100	1	1
101	1	1
110	1	1
111	0	0

Review: Minterms & Maxterms

ABC	Minterm	name	Maxterm	name
000	$A' \cdot B' \cdot C'$	m_0	$A+B+C$	M_0
001	$A' \cdot B' \cdot C$	m_1	$A+B+C'$	M_1
010	$A' \cdot B \cdot C'$	m_2	$A+B'+C$	M_2
011	$A' \cdot B \cdot C$	m_3	$A+B'+C'$	M_3
100	$A \cdot B' \cdot C'$	m_4	$A'+B+C$	M_4
101	$A \cdot B' \cdot C$	m_5	$A'+B+C'$	M_5
110	$A \cdot B \cdot C'$	m_6	$A'+B'+C$	M_6
111	$A \cdot B \cdot C$	m_7	$A'+B'+C'$	M_7

$(m_i)' = M_i$ according to De Morgan

$m_i = 1 \iff$ input row i is selected

$M_i = 1 \iff$ input row i is NOT selected

Review: Canonical Representations

- Problem: $F = 1$ if and only if $A \cdot B = (B + C)'$
- Step 1: Lay out the truth table
- Step 2: Derive canonical forms

– Canonical sum ($F=1$ if input “hits” the on-set)

$$\begin{aligned}
 F &= A' \cdot B' \cdot C + A \cdot B' \cdot C + A \cdot B \cdot C' + A \cdot B \cdot C \\
 &= \sum_{A,B,C} (1,5,6,7)
 \end{aligned}$$

– Canonical product ($F=1$ if input “avoids” the off-set)

$$\begin{aligned}
 F &= (A' \cdot B' \cdot C')' (A' \cdot B \cdot C')' (A' \cdot B \cdot C)' (A \cdot B' \cdot C')' \\
 &= (A + B + C)(A + B' + C)(A + B' + C')(A' + B + C) \\
 &= \prod_{A,B,C} (0,2,3,4)
 \end{aligned}$$

ABC	F
000	0
001	1
010	0
011	0
100	0
101	1
110	1
111	1

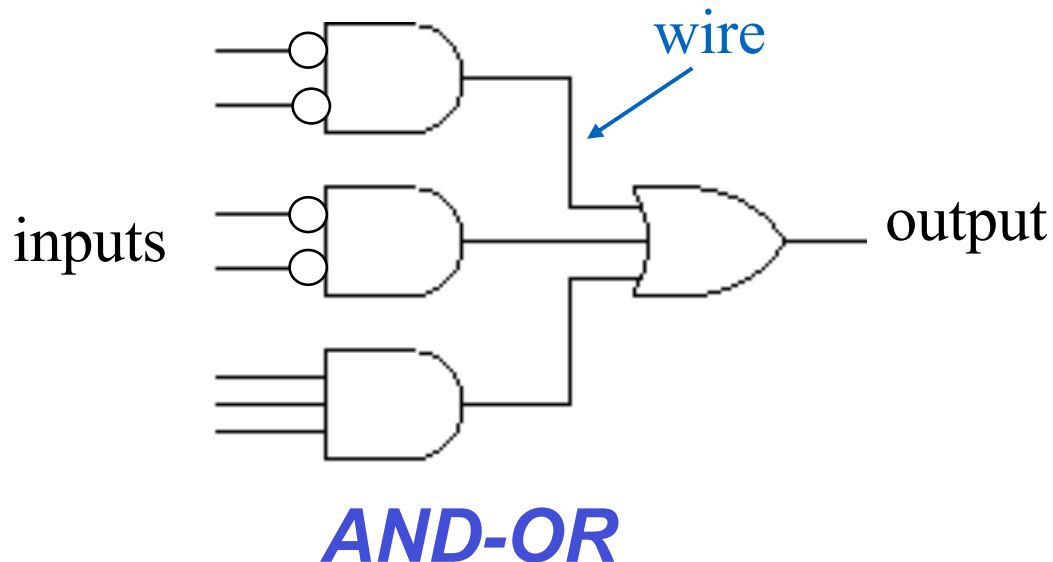
- Step 3: Simplification (this lecture)

Combinational Logic

- **Outputs depend *only* on current inputs**
 - **Example: Detect if the input is an odd number**
 - **Can be represented in two-level or multi-level forms**
- **In contrast, sequential logic has “memory” or “state”**
 - **Example: Detect if the last two inputs are odd**
 - **We’ll cover sequential logic later**

Two-Level Logic: Sum-of-Products

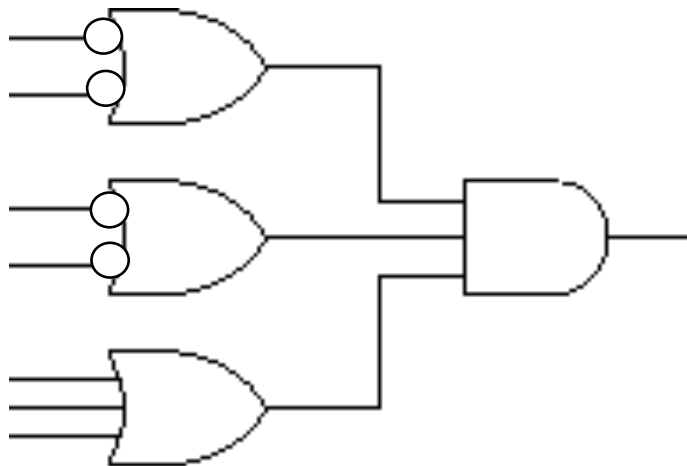
- **Sum of product terms (SOP)**
 - e.g., $A' \cdot B' + A' \cdot C' + A \cdot B \cdot C$
- **Circuits look something like this**
 - *A bubble* indicates the signal is inverted



Two-Level Logic: Product-of-Sums

- **Product of sum terms (POS)**
 - e.g., $(A'+C') \cdot (B'+C') \cdot (A+B+C)$

- **Circuits look something like this**



OR-AND

Combinational logic can be expressed as SOP or POS

Algebraic Simplification

- Apply theorems to canonical sum (or product) to reduce (1) the number of terms, and (2) the number of literals in each term
- Results in a more compact expression and lower cost digital logic implementation
- We focus on *minimizing two-level logic* in this lecture

Algebraic Simplification Example

- **Binary adder**

- inputs: A, B, Carry-in (Cin)

- outputs: Sum, Carry-out (Cout)



- **Truth Table → Canonical sum**

A	B	C _{in}	C _{out}	S
0	0	0	0	0
0	0	1	0	1
0	1	0	0	1
0	1	1	1	0
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	1	1

$$S = A' \cdot B' \cdot C_{in} + A' \cdot B \cdot C_{in}' + A \cdot B' \cdot C_{in}' + A \cdot B \cdot C_{in}$$

$$C_{out} = A' \cdot B \cdot C_{in} + A \cdot B' \cdot C_{in} + A \cdot B \cdot C_{in}' + A \cdot B \cdot C_{in}$$

Algebraic Simplification Example

$$\begin{aligned} \text{Cout} &= A' \cdot B \cdot \text{Cin} + A \cdot B' \cdot \text{Cin} + A \cdot B \cdot \text{Cin}' + A \cdot B \cdot \text{Cin} \\ &= A' \cdot B \cdot \text{Cin} + A \cdot B' \cdot \text{Cin} + A \cdot B \cdot \text{Cin}' + A \cdot B \cdot \text{Cin} + A \cdot B \cdot \text{Cin} && \text{(idempotency)} \\ &= B \cdot \text{Cin} + A \cdot B' \cdot \text{Cin} + A \cdot B \cdot \text{Cin}' + A \cdot B \cdot \text{Cin} && \text{(combining)} \\ &= B \cdot \text{Cin} + A \cdot B' \cdot \text{Cin} + A \cdot B \cdot \text{Cin}' + A \cdot B \cdot \text{Cin} + A \cdot B \cdot \text{Cin} && \text{(idempotency)} \\ &= B \cdot \text{Cin} + A \cdot \text{Cin} + A \cdot B \cdot \text{Cin}' + A \cdot B \cdot \text{Cin} && \text{(combining)} \\ &= B \cdot \text{Cin} + A \cdot \text{Cin} + A \cdot B && \text{(combining)} \end{aligned}$$

We can apply these theorems in a more intuitive fashion using a *Karnaugh Map*

Reduction in Hardware Cost

$$\begin{aligned}C_{out} &= A' \cdot B \cdot C_{in} + A \cdot B' \cdot C_{in} + A \cdot B \cdot C_{in}' + A \cdot B \cdot C_{in} \\ &= B \cdot C_{in} + A \cdot C_{in} + A \cdot B\end{aligned}$$

3 inverters

4 three-input ANDs

1 four-input OR



3 two-input ANDs

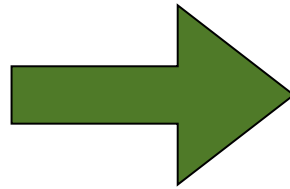
1 three-input OR

Karnaugh Map (K-Map)

- **Idea: Use combining and idempotency theorems *visually* to simplify canonical forms into two-level SOP or POS**
- **Multidimensional representation of a truth table**
- **Adjacent cells represent minterms (or maxterms) that differ by exactly one literal**
 - **Cyclic encoding along each dimension**
- **At most two variables per dimension**

K-Map for Cout

A	B	Cin	S	Cout
0	0	0	0	0
0	0	1	1	0
0	1	0	1	0
0	1	1	0	1
1	0	0	1	0
1	0	1	0	1
1	1	0	0	1
1	1	1	1	1



		AB			
		00	01	11	10
Cin	0	0	0	1	0
	1	0	1	1	1

Some K-Map Definitions

- 1-cell: Minterm of a canonical sum
- Implicant: Set of adjacent 1-cells
 - A rectangle in K-Map
 - Number of cells contained must be a power of 2
- Prime implicant: Implicant that cannot be contained in a larger implicant

AB	00	01	11	10
Cin 0	0	0	1	0
Cin 1	0	1	1	1

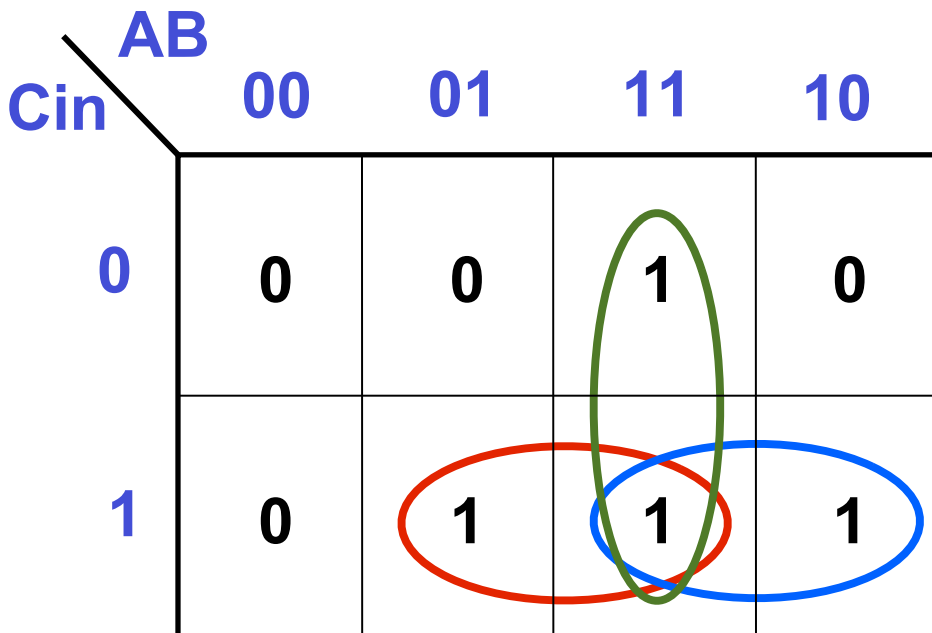
A prime implicant

SOP Minimization Using K-Map

- **Goal:** Cover all 1-cells with the minimum number of prime implicants
- **Procedure**
 - Plot 1's corresponding to minterms of function
 - Circle largest possible rectangular sets of 1's
 - Must be power of 2
 - Including “wrap-around” sets
 - Repeat until all minterms are covered
 - OR product terms derived from each circle
- **Minimizes the gate count and inputs in the SOP form**
 - Solution may not be unique

Simplifying Cout Using a K-Map

$$\begin{aligned}
 \text{Cout} &= A' \cdot B \cdot \text{Cin} + A \cdot B' \cdot \text{Cin} + A \cdot B \cdot \text{Cin}' + A \cdot B \cdot \text{Cin} \\
 &= A' \cdot B \cdot \text{Cin} + A \cdot B' \cdot \text{Cin} + A \cdot B \cdot \text{Cin}' + \mathbf{A \cdot B \cdot \text{Cin}} + \mathbf{A \cdot B \cdot \text{Cin}} && \text{(idempotency)} \\
 &= \mathbf{B \cdot \text{Cin}} + A \cdot B' \cdot \text{Cin} + A \cdot B \cdot \text{Cin}' + A \cdot B \cdot \text{Cin} && \text{(combining)} \\
 &= B \cdot \text{Cin} + A \cdot B' \cdot \text{Cin} + A \cdot B \cdot \text{Cin}' + \mathbf{A \cdot B \cdot \text{Cin}} + \mathbf{A \cdot B \cdot \text{Cin}} && \text{(idempotency)} \\
 &= B \cdot \text{Cin} + \mathbf{A \cdot \text{Cin}} + A \cdot B \cdot \text{Cin}' + A \cdot B \cdot \text{Cin} && \text{(combining)} \\
 &= B \cdot \text{Cin} + A \cdot \text{Cin} + \mathbf{A \cdot B} && \text{(combining)}
 \end{aligned}$$



Circles (ovals) indicate applying combining theorem

Idempotency theorem allows circles to overlap

3 Variable K-Map Example

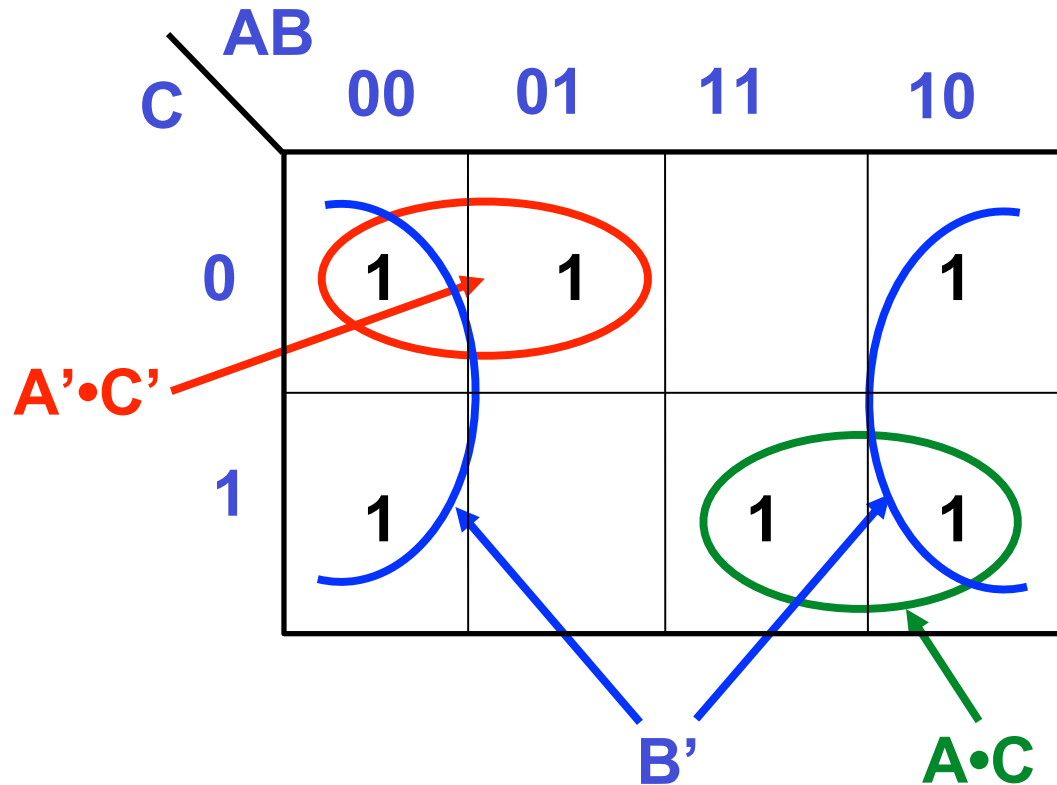
$$F = \sum_{A,B,C}(4,5,6,7)$$

		AB			
		00	01	11	10
C	0			1	1
	1			1	1

$$F = A$$

Another Example

$$F = \sum_{A,B,C} (0, 1, 2, 4, 5, 7)$$



$$F = A' \cdot C' + B' + A \cdot C$$

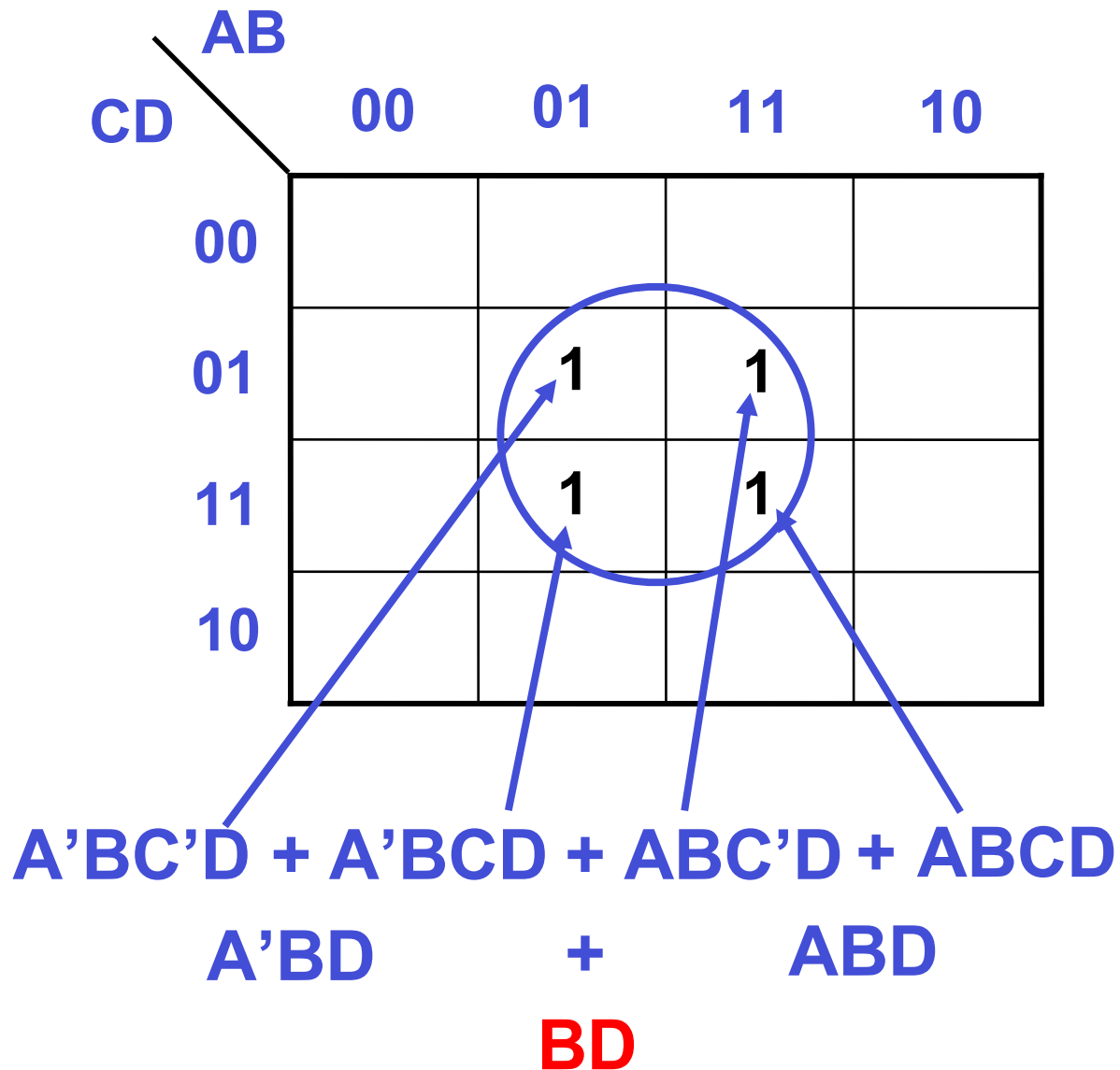
Intuition Behind K-Map

- **Spatial encoding!**

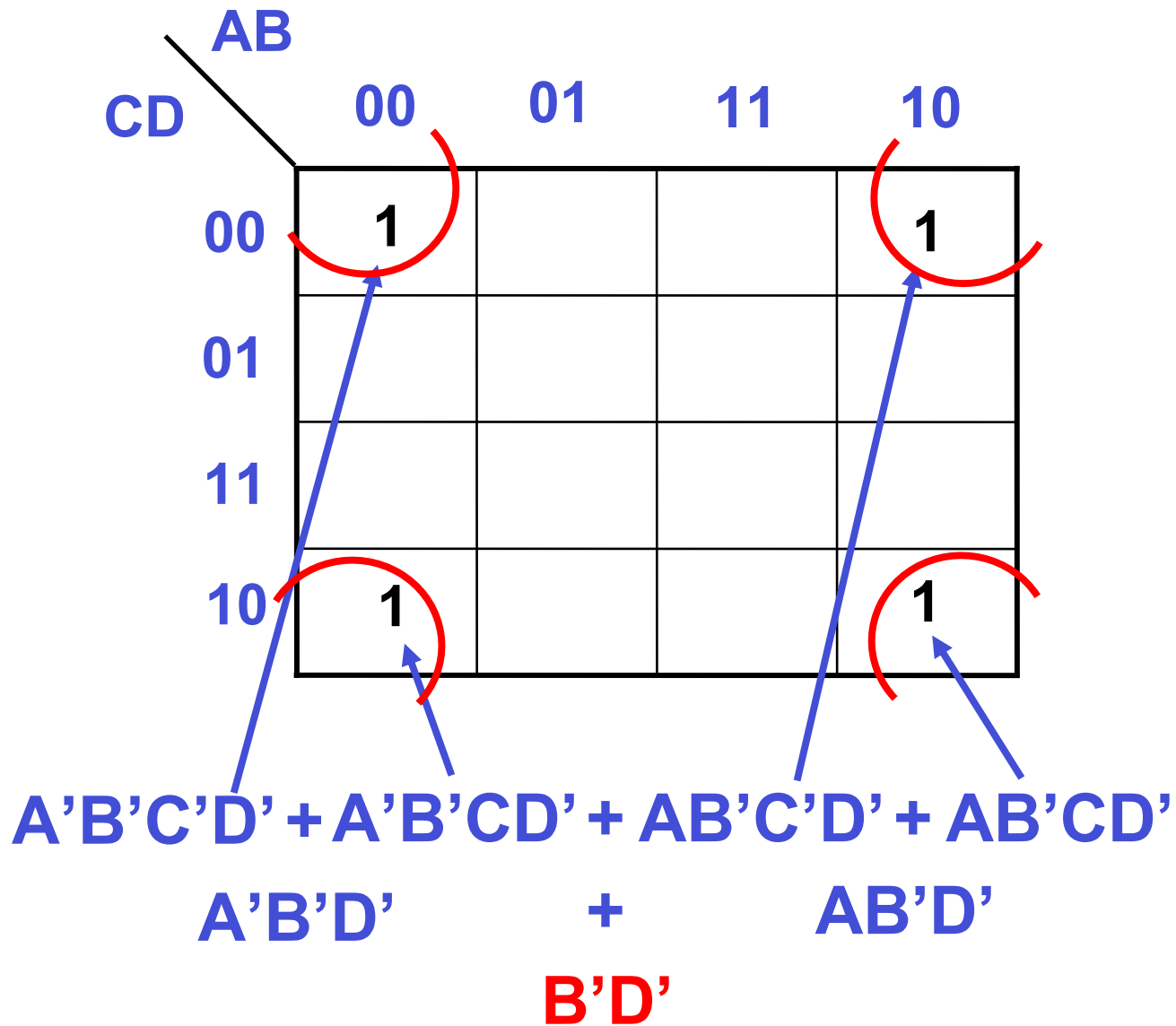
4 Variable K-Map

		AB			
		00	01	11	10
CD	00				
	01				
	11				
	10				

Combining Theorem in Action

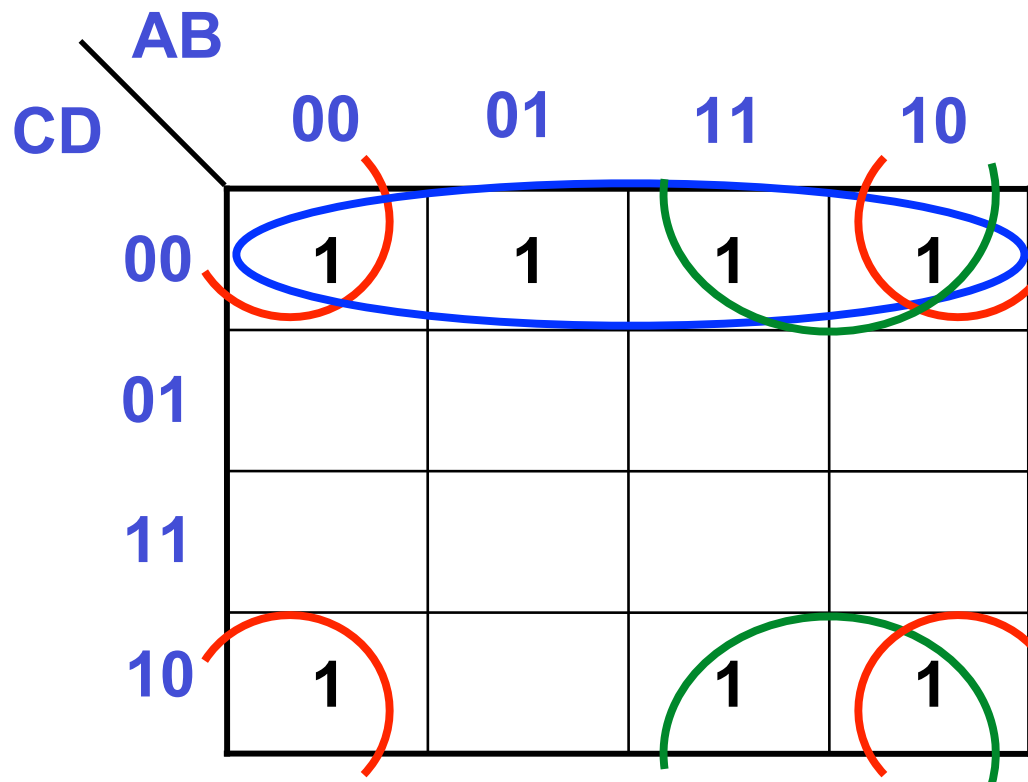


Combining Theorem in Action



4 Variable Karnaugh Map Example

- Detect all even digits from 0 to 15 except 6
 - $F = \sum_{A,B,C,D}(0,2,4,8,10,12,14)$



$$F = B'D' + AD' + C'D'$$

Minimizing Product-of-Sums

- **Procedure**
 - Plot 0's corresponding to maxterms of function
 - Circle largest possible rectangular sets of 0's
 - Must be power of 2
 - Including “wrap-around” sets
 - Repeat until all maxterms are covered
 - AND sum terms derived from each circle

Product-of-Sums Example

$$F = \prod_{A,B,C,D} (0, 2, 6, 7, 8, 10)$$

		AB			
		00	01	11	10
CD	00	0			0
	01				
	11		0		
	10	0	0		0

Corners:

$$B + D$$

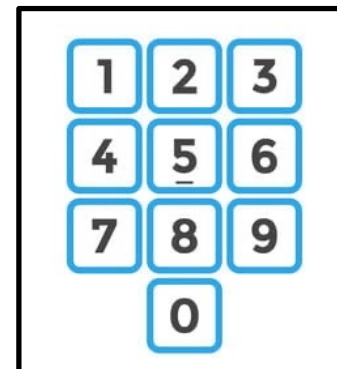
Other:

$$A + B' + C'$$

$$F = (B + D) \cdot (A + B' + C')$$

Don't Cares Combinations

- Sometimes, for a given input combination, the output is unspecified or irrelevant
 - Such as an input combination that will never occur based on the problem definition
 - We call this input combination a Don't Care
 - Represent as a 'd' (or 'x') in the truth table and K-map
- Example: Detect when an odd decimal digit (except 3) is pressed on a keypad
 - Four input bits are still required, but inputs 0-9 would appear
 - 10-15 are “don't care” values



Don't Cares in Karnaugh Map

- Don't cares can be used as 1-cells as needed for minimizing SOP or 0-cells for POS
- Only circle if doing so creates a larger prime implicant (and thus a more minimal expression)

- Example: Detect when an odd decimal digit (except 3) is pressed on a keypad
 - Inputs 10-15 will never occur

		AB			
		00	01	11	10
CD	00	0	0	d	0
	01	1	1	d	1
	11	0	1	d	d
	10	0	0	d	d

$$F = BD + C'D$$

Next Class

Combinational Building Blocks (H&H 2.8)