Towards a Deeper Understanding of Training Quantized Neural Networks

Hao Li *1 Soham De *1 Zheng Xu 1 Christoph Studer 2 Hanan Samet 1 Tom Goldstein 1

Abstract

Training neural networks with coarsely quantized weights is a key step towards learning on embedded platforms that have limited computing resources, memory capacity, and power consumption. Numerous recent publications have studied methods for training quantized networks, but these studies have been purely experimental. In this work, we investigate the theory of training quantized neural networks by analyzing the convergence properties of some commonly used methods. Our main result shows that training algorithms that exploit high-precision representations have an important annealing property that purely quantized training methods lack, which explains many of the observed empirical differences between these types of algorithms.

1. Introduction

Neural networks (NNs) are an integral part of numerous state-of-the-art computer vision and natural language processing tasks. Because of their high memory requirements and computational complexity, networks are usually trained using powerful hardware. There is an increasing interest in training and deploying neural networks directly on battery-powered devices, such as cell phones or other platforms. Such low-power embedded systems are, however, memory and power limited, and in some cases lack basic support for floating-point arithmetic.

To make neural nets practical on embedded systems, many researchers have focused on training nets with coarsely quantized weights. For example, weights may be constrained to integer/binary values, or represented using low-precision (8 bits or less) fixed-point numbers. Quantized nets offer the potential of superior memory and computation efficiency, while achieving performance competitive with state-of-the-art high-precision nets. Quantized weights can dramatically reduce memory size and access bandwidth, increase power efficiency, exploit hardware-friendly bitwise operations, and accelerate inference throughput (Marchesi et al., 1993; Courbariaux et al., 2016; Rastegari et al., 2016).

Handling low-precision weights is difficult and motivates interest in new training methods. When learning rates are small, stochastic gradient methods make small updates to the weights. Binarization/discretization of weights after each training iteration “rounds off” these small updates and causes training to stagnate (Courbariaux et al., 2016). Thus, the naive approach of quantizing weights using rounding yields poor results when weights are represented using a small number of bits. Other approaches include classical stochastic rounding methods (Gupta et al., 2015), as well as schemes that combine full-precision floating-point weights with discrete rounding procedures (Courbariaux et al., 2015). While some of these schemes seem to work in practice, results in this area are largely experimental, and little work has been devoted to explaining the excellent performance of some methods, the poor performance of others, and the important differences in behavior between these methods.

Contributions This paper studies quantized training methods from a theoretical perspective, with the goal of understanding the differences in behavior, and reasons for success or failure, of various methods. In particular, we show that recent powerful methods like BinaryConnect (BC) (Courbariaux et al., 2015) are capable of concentrating on minimizers and solving discrete problems up to a high level of accuracy. On the other hand, we show that classical stochastic rounding (SR) methods (Gupta et al., 2015) lack an important annealing property that is required for non-convex optimization. This explains the large differences observed in empirical performance between classical SR methods and the more powerful BC methods.

2. Related Work

The arithmetic operations of deep networks can be encoded down to 8-bit fixed-point without significant deterioration

---

*Equal contribution. Author ordering determined by coin flip.
1Department of Computer Science, University of Maryland, College Park, MD
2Department of Electrical and Computer Engineering, Cornell University, Ithaca, NY. Correspondence to: Hao Li <haoli@cs.umd.edu>, Soham De <sohamde@cs.umd.edu>.

Presented at the ICML 2017 Workshop on Principled Approaches to Deep Learning, Sydney, Australia, 2017. Copyright 2017 by the author(s).
inference performance (Gupta et al., 2015; Lin et al., 2016a; Hwang & Sung, 2014; Lin et al., 2016b; Li et al., 2016). The most extreme scenario of quantization is binarization, in which only 1-bit (two states) is available for weight representation (Kim & Smaragdis, 2015; Courbariaux et al., 2015; 2016; Rastegari et al., 2016; Hubara et al., 2016; Baldassi et al., 2015).

Previous work on obtaining a quantized neural network (NN) can be divided into two categories: quantizing pre-trained models with or without retraining (Hwang & Sung, 2014; Anwar et al., 2015; Lin et al., 2016a; Zhu et al., 2017; Zhou et al., 2017), and training a quantized model from scratch (Gupta et al., 2015; Courbariaux et al., 2015; Rastegari et al., 2016; Courbariaux et al., 2016; Zhou et al., 2016). We focus on approaches that belong to the second category, as they can be used for both training and inference under constrained resources.

For training quantized NNs from scratch, many authors suggest maintaining a high-precision copy of the weights while feeding quantized weights into backprop (Courbariaux et al., 2015; Hubara et al., 2016; Rastegari et al., 2016; Zhou et al., 2016), which results in good empirical performance. There are limitations in using such methods on low-power devices, however, where floating-point arithmetic is not available or not desirable. Another widely used solution using only low-precision weights is stochastic rounding (Höhfeld & Fahlman, 1992; Gupta et al., 2015). Experiments show that networks using 16-bit fixed-point representations with stochastic rounding can deliver results nearly identical to 32-bit floating-point computations (Gupta et al., 2015), while lowering the precision down to 3-bit fixed-point often results in a significant performance degradation (Miyashita et al., 2016). Bayesian learning has also been applied to train binary networks (Soudry et al., 2014; Cheng et al., 2015). A more comprehensive review can be found in (Rastegari et al., 2016).

3. Training Quantized Neural Nets

We consider problems of the form:

$$\min_{w \in \mathbb{W}} F(w) := \frac{1}{m} \sum_{i=1}^{m} f_i(w),$$

where the objective function decomposes into a sum over many functions $f_i : \mathbb{R}^d \rightarrow \mathbb{R}$. Neural networks have objective functions of this form where each $f_i$ is a non-convex loss function. When floating-point representations are available, the standard method for training neural networks is stochastic gradient descent (SGD), which on each iteration selects a function $f$ randomly from $\{f_1, f_2, \ldots, f_m\}$, and then computes

SGD: $w^{t+1} = w^t - \alpha_t \nabla f(w^t),$ \hspace{1cm} (2)

for some learning rate $\alpha_t$. In this paper, we consider the problem of training convolutional neural networks (CNNs). In CNNs, each convolutional layer $i$ accepts inputs $x_i$ and generates the response corresponding to filter $j$ with linear transformation $a_{i,j} = w_{i,j} \ast x_i + b_j$, where $w_{i,j}$ represent the vector of weights of filter $j$ at layer $i$, and $\ast$ denotes the convolution operation. The neuron output is generated by applying a non-linear activation function $h(\cdot)$ to the filter response, i.e., $x_{i+1} = h(a_{i})$, which is used as the input to the next layer. Convolutions are expensive; low precision weights can be used to accelerate the convolutions by replacing expensive multiplications with binary operations or dedicated high-performance computer arithmetic.

To train neural networks using a low-precision representation of the weights, a quantization function $Q(\cdot)$ is required that converts a real-valued number $w$ into a quantized version $\hat{w} = Q(w)$. We use the same notation for quantizing vectors, where we assume $Q$ acts on each dimension of the vector. Different quantized optimization routines can be defined by selecting different quantizers, and also by selecting when quantization happens during optimization. The common options are:

**Deterministic Rounding (R)** A basic uniform or deterministic quantization function converts a floating point value to the closest quantized value as:

$$Q_d(w) = \text{sign}(w) \cdot \Delta \cdot \left\lfloor \frac{|w|}{\Delta} + \frac{1}{2} \right\rfloor,$$ \hspace{1cm} (3)

where $\Delta$ denotes the quantization step or resolution, i.e., the smallest positive number that is representable. One exception to this definition is when we consider binary weights, where all weights are constrained to have two values $w \in \{-1, 1\}$ and uniform rounding becomes $Q_d(w) = \text{sign}(w)$.

The deterministic rounding SGD maintains quantized weights with updates of the form:

Deterministic Rounding: $w_{b}^{t+1} = Q_d(w_{b}^t - \alpha_t \nabla \hat{f}(w_{b}^t)),$ \hspace{1cm} (4)

where $w_{b}$ denotes the low-precision weights, which are quantized using $Q_d$ immediately after applying the gradient descent update. If gradient updates are significantly smaller than the quantization step, this method loses gradient information and weights may never be modified from their starting values.

**Stochastic Rounding (SR)** The quantization function for stochastic rounding is defined as:

$$Q_s(w) = \Delta \cdot \begin{cases} \left\lfloor \frac{|w|}{\Delta} \right\rfloor + 1 & \text{for } p \leq \frac{|w|}{\Delta} < \left\lfloor \frac{|w|}{\Delta} \right\rfloor, \\ \left\lfloor \frac{|w|}{\Delta} \right\rfloor & \text{otherwise}, \end{cases}$$ \hspace{1cm} (5)
where \( p \in [0, 1] \) is produced by a uniform random number generator. This operator is non-deterministic, and rounds its argument up with probability \( w/\Delta - \lfloor w/\Delta \rfloor \), and down otherwise. This quantizer satisfies the important property:

\[
\mathbb{E}[Q_x(w)] = w.
\]

Similar to the deterministic rounding method, the SR optimization method also maintains quantized weights with updates of the form:

**Stochastic Rounding:** \( w_b^{t+1} = Q_b(w_b^t - \alpha_t \nabla f(w_b^t)) \). (6)

**BinaryConnect (BC)** The BinaryConnect algorithm (Courbariaux et al., 2015) accumulates gradient updates using a full-precision buffer \( w_r \), and quantizes weights only just before gradient computations. BinaryConnect uses updates of the form:

\[
\text{BinaryConnect: } w_r^{t+1} = w_r^t - \alpha_t \nabla \tilde{f}(Q(w_r^t)). \tag{7}
\]

Either stochastic rounding \( Q_s \) or deterministic rounding \( Q_d \) can be used for quantizing the weights \( w_r \), but in practice, \( Q_d \) is the common choice. The original BinaryConnect paper constrains the weights to be \( \{-1, 1\} \), which can be generalized to \( \{-\Delta, \Delta\} \). A more recent method, BinaryWeights-Net (BWN) (Rastegari et al., 2016), allows different filters to have different scales for quantization, which often results in better performance on large datasets.

**Notation** For the rest of the paper, we use \( Q \) to denote both \( Q_s \) and \( Q_d \) unless the situation requires this to be distinguished. We also drop the subscripts on \( w_r \) and \( w_b \), and simply write \( w \).

## 4. Analyzing the BinaryConnect (BC) Method

Here we present convergence guarantees for the BinaryConnect (BC) algorithm, with updates of the form (7), and we show that BinaryConnect has the property of concentrating its iterates on minimizers over time. To do this, we make the following assumptions that are standard for analyzing convergence rates of optimization algorithms over convex functions.

**Assumption 1.** We assume that \( F \) is \( \mu \)-strongly convex, i.e., it satisfies the following inequality:

\[
\langle \nabla F(w'), w - w' \rangle \leq F(w) - F(w') - \frac{\mu}{2} \|w - w'\|^2.
\]

**Assumption 2.** We assume that the magnitude of each \( \nabla f_i \) is uniformly upper-bounded as:

\[
\|\nabla f_i(w)\|^2 \leq G^2, \quad \forall w \in W, \forall i = 1, 2, \ldots, m.
\]

**Assumption 3.** We assume that the domain is bounded as:

\[
\|w - w^*\| \leq D, \quad \forall w \in W.
\]

In this case, the rounding algorithm clips values that leave the domain. For example, in the binary case, the rounding algorithm never returns a value outside of \( \{-1, 1\} \).

In addition to these standard assumptions, we also assume the following.

**Assumption 4.** We assume that the Hessian satisfies

\[
\|\nabla^2 F_i(x) - \nabla^2 F_i(y)\| \leq L_2 \|x - y\|,
\]

for some constant \( L_2 > 0 \). Equivalently, this translates to the condition:

\[
\nabla F_i(y) = \nabla F_i(x) + \nabla^2 F_i(x)(y - x) + \delta,
\]

where

\[
\|\delta\| \leq \frac{L_2}{2} \|x - y\|^2.
\]

While this is a slightly non-standard assumption, we will see below that it enables us to gain better insights into the behavior of the BinaryConnect algorithm. We now analyze the BinaryConnect algorithm.

**Theorem 1.** Given Assumptions 1-4 and learning rates given by \( \alpha_t = \frac{1}{m(t+1)} \), the BinaryConnect algorithm, with updates of the form (7) using stochastic rounding quantization function \( Q_s \), converges as:

\[
\mathbb{E}[F(\bar{w}^T) - F(w^*)] \leq \frac{(1 + \log(T+1))G^2}{2\mu T} + \frac{DL_2 \sqrt{d} \Delta}{2T},
\]

where \( \bar{w}^T \) denotes the iterate average: \( \bar{w}^T = \frac{1}{T} \sum_{t=1}^T w_t \).

**Proof.** From the update rule (7), we get

\[
w^{t+1} = w^t - \alpha_t \nabla \tilde{f}(Q(w^t))
\]

\[
= w^t - \alpha_t \nabla \tilde{f}(w^t + r^t)
\]

\[
= w^t - \alpha_t [\nabla \tilde{f}(w^t) + \nabla^2 \tilde{f}(w^t) r^t + \tilde{r}^t],
\]

where \( \|\tilde{r}^t\| \leq \frac{L_2}{2} \|r^t\|^2 \) from Assumption 4. Note that in general \( r^t \) has mean zero while \( \tilde{r}^t \) does not.

Subtracting by the optimal \( w^* \), taking norm, and taking expectation conditioned on \( w^t \), we get:

\[
\mathbb{E}[\|w^{t+1} - w^*\|^2]
\]

\[
= \mathbb{E}[\|w^t - w^*\|^2 - 2\alpha_t \mathbb{E}\langle w^t - w^*, \nabla \tilde{f}(w^t + r^t) \rangle + \alpha_t^2 \mathbb{E}[\|\nabla \tilde{f}(w^t + r^t)\|^2]
\]

\[
\leq \mathbb{E}[\|w^t - w^*\|^2 - 2\alpha_t \mathbb{E}\langle w^t - w^*, \nabla F(w^t) + \tilde{r}^t \rangle + \alpha_t^2 \mathbb{E}[\|\nabla \tilde{f}(w^t + r^t)\|^2]
\]

\[
= \mathbb{E}[\|w^t - w^*\|^2 - 2\alpha_t \mathbb{E}\langle w^t - w^*, \nabla F(w^t) \rangle + \alpha_t^2 G^2 - 2\alpha_t \mathbb{E}\langle w^t - w^*, \tilde{r}^t \rangle],
\]

where the inequality uses Assumptions 2 and 4. Using Assumption 3 and observing that the quantization error for
Towards a Deeper Understanding of Training Quantized Neural Networks

BC-SGD can always be upper-bounded as \( \|r^t\| \leq \sqrt{2} \Delta \), we get:

\[
-2\alpha_t E(w^t - w^*)^2 \leq 2\alpha_t D\|w^t\|^2 \\
\leq 2\alpha_t D L_2 \|r^t\|^2 \\
\leq \alpha_t D L_2 \sqrt{2} \Delta.
\]

Thus, using Assumption 1, we get:

\[
E\|w^{t+1} - w^*\|^2 \leq (1 - \alpha_t^2)\|w^t - w^*\|^2 + \alpha_t D L_2 \sqrt{2} \Delta \\
- 2\alpha_t (F(w^t) - F(w^*)) + \alpha_t^2 G^2.
\]

Re-arranging the terms, and taking expectation we get:

\[
E(F(w^t) - F(w^*)) \leq \left( \frac{1}{2\alpha_t} - \frac{\mu}{2} \right) E\|w^t - w^*\|^2 + \frac{\alpha_t G^2}{2} \\
- \frac{1}{2\alpha_t} E\|w^t + 1 - w^*\|^2 + \frac{D L_2 \sqrt{2} \Delta}{2}.
\]

Assume that the step sizes decrease with the rate \( \alpha_t = 1/\mu(t+1) \). Then we have:

\[
E(F(w^t) - F(w^*)) \leq \frac{G^2}{2\mu(t+1)} + \frac{D L_2 \sqrt{2} \Delta}{2} \\
- \frac{\mu(t+1)}{2} E\|w^t - w^*\|^2 + \frac{D L_2 \sqrt{2} \Delta}{2}.
\]

Averaging over \( t = 0 \) to \( T \), we get a telescoping sum on the right hand side, which yields:

\[
\frac{1}{T} \sum_{t=0}^{T} E(F(w^t) - F(w^*)) \\
\leq \frac{G^2}{2\mu T} \sum_{t=0}^{T} \frac{1}{t+1} + \frac{D L_2 \sqrt{2} \Delta}{2} - \frac{\mu(t+1)}{2} E\|w^T - w^*\|^2 + \frac{D L_2 \sqrt{2} \Delta}{2}.
\]

Using Jensen’s inequality, we have:

\[
E(F(\bar{w}) - F(w^*)) \leq \frac{1}{T} \sum_{t=0}^{T} E(F(w^t) - F(w^*)),
\]

where \( \bar{w} = \frac{1}{T} \sum_{t=0}^{T} w^t \), the average of the iterates. The final convergence result follows.

\[\Box\]

Thus, we see that BinaryConnect converges until it reaches an “accuracy floor”, which is determined by the quantization error \( \Delta \) and \( L_2 \). Now, consider a quadratic least-squares problem, \( F(w) = \frac{1}{2} \|Aw - b\|^2 \). Here, the gradient of \( F \) is linear: \( \nabla F(w) = Aw - b \), and the Hessian is constant. Thus, \( L_2 = 0 \) and we get the following corollary.

\textbf{Corollary 1.} Assume that \( F \) is quadratic of the form \( F(w) = \frac{1}{2} \|Aw - b\|^2 \) and the learning rates are given by \( \alpha_t = \frac{1}{\mu(t+1)} \). Given Assumptions 1-4, the BC algorithm with updates of the form (7) and using the stochastic rounding quantization function \( Q_s \), yields

\[
E[F(\bar{w}) - F(w^*)] \leq \frac{(1 + \log(T + 1)G^2)}{2\mu T}.
\]

We see that accumulating the real-valued weights in BC allows it to converge to the \textit{true minimizer} of quadratic losses. Furthermore, this suggests that, when the function behaves like a quadratic on the distance scale \( \Delta \), one would expect BC to effectively concentrate on minimizers.

\section{5. Analyzing the Stochastic Rounding Method}

We can rewrite the update rule (6) of the Stochastic Rounding (SR) algorithm as:

\[
w^{t+1} = w^t - \alpha_t \nabla \tilde{f}(w^t) + r^t,
\]

where \( r^t = Q_s(w^t - \alpha_t \nabla \tilde{f}(w^t)) - w^t + \alpha_t \nabla \tilde{f}(w^t) \) denotes the quantization error on the \( t \)-th iteration.

\subsection{5.1. Bound on Quantization Error}

We want to bound the error \( r^t \) in expectation. To this end, we present the following lemma.

\textbf{Lemma 1.} The error \( r^t \) on each iteration \( t \) of the Stochastic Rounding algorithm with updates given by (8) can be bounded, in expectation, as:

\[
E\|r^t\|^2 \leq \sqrt{2} \Delta \alpha_t G.
\]

\textbf{Proof.} We want to bound the quantization error \( r^t \). Consider the \( i \)-th entry in \( r^t \) denoted by \( r_i^t \). Similarly, we define \( w_i^t \) and \( \nabla \tilde{f}(w_i^t) \). Choose some random number \( p \in [0, 1] \). The stochastic rounding operation produces a value of \( r_i^t \) given by

\[
r_i^t = Q_s(w_i^t - \alpha_t \nabla \tilde{f}(w_i^t)) - w_i^t + \alpha_t \nabla \tilde{f}(w_i^t)
\]

\[
= \Delta \cdot \begin{cases}
\frac{\alpha_t \nabla \tilde{f}(w_i^t)}{\Delta} + \frac{-\alpha_t \nabla \tilde{f}(w_i^t)}{\Delta} + 1, & \text{for } p \leq q,
\frac{\alpha_t \nabla \tilde{f}(w_i^t)}{\Delta} + \frac{-\alpha_t \nabla \tilde{f}(w_i^t)}{\Delta}, & \text{otherwise},
\end{cases}
\]

\[
= \Delta \cdot \begin{cases}
-q + 1, & \text{for } p \leq q,
-q, & \text{otherwise}.
\end{cases}
\]

where we write \( q = \frac{\alpha_t \nabla \tilde{f}(w_i^t)}{\Delta} - \frac{-\alpha_t \nabla \tilde{f}(w_i^t)}{\Delta} \) and \( q \in [0, 1] \).
Now we have
\[
E_p[\|r^t\|^2] \leq \Delta^2((-q + 1)^2q + (-q)^2(1-q))
\]
= \Delta^2q(1-q)
\]
\[\leq \Delta^2 \min\{q, 1-q\}.
\]
Because \(\min\{q, 1-q\} \leq \left| \alpha_t \nabla \tilde{f}(w^t)_i \right| / \Delta \), it follows that
\[
E_p[\|r^t\|^2] \leq \Delta^2 \left| \frac{\alpha_t \nabla \tilde{f}(w^t)_i}{\Delta} \right| \leq \Delta \left| \alpha_t \nabla \tilde{f}(w^t)_i \right|.
\]
Summing over the index \(i\) yields
\[
E_p[\|r^t\|^2] \leq \Delta \alpha_t \|\nabla \tilde{f}(w^t)\|_1 \leq \sqrt{d} \alpha_t \|\nabla \tilde{f}(w^t)\|_2. \tag{9}
\]
Now, \((E[\|\nabla \tilde{f}(w^t)\|_2^2] \leq \|\nabla \tilde{f}(w^t)\|_2^2 \leq G^2\). Plugging this into (9) yields
\[
E_p[\|r^t\|^2] \leq \sqrt{d} \Delta \alpha_t G. \tag{10}
\]
\[
\square
\]

From Lemma 1, we see that the rounding error per step decreases as the learning rate \(\alpha_t\) decreases. This is intuitive since the probability of an entry in \(w^{t+1}\) differing from \(w^t\) is small when the gradient update is small relative to \(\Delta\).

To explain the large observed differences in asymptotic behavior between the BinaryConnect and Stochastic Rounding algorithms, we now study the asymptotic, long-term behavior of the SR algorithm as the learning rate gets small.

Note that Section 5 makes no convexity assumptions, and the proposed theoretical results are directly applicable to neural networks.

5.2. Asymptotic Behavior of Stochastic Rounding

Typical (continuous-valued) SGD methods have an important exploration-exploitation tradeoff: when the learning rate is large, the algorithm explores by moving quickly between states. Exploitation happens when the learning rate is small. In this case, noise averaging causes the algorithm to behave more like deterministic gradient descent, where the algorithm more greedily pursues local minimizers with lower loss function values. Thus, the distribution of iterates produced by the algorithm becomes increasingly concentrated near minimizers as the learning rate vanishes (see, e.g., the large-deviations estimates in (Lan et al., 2012)).

BC maintains this property as well—indeed, we saw in Corollary 1 a class of problems for which the iterates concentrate near the minimizer for small \(\alpha_t\).

Here we show that the SR method lacks this important tradeoff: as the stepsize gets small and the algorithm slows down, the quality of the iterates produced by the algorithm does not improve, and the algorithm does not become progressively more likely to produce low-loss iterates.

To understand this problem conceptually, consider the simplified case of a single-variable optimization problem starting at \(x^0 = 0\) with \(\Delta = 1\), as depicted in Figure 1. On each iteration, the algorithm computes a stochastic approximation \(\nabla \tilde{f}\) of the gradient by sampling from a distribution, which we call \(p\). This gradient is then multiplied by the stepsize to get \(\alpha \nabla \tilde{f}\). The probability of moving to the right (or left) is then roughly proportional to the magnitude of \(\alpha \nabla \tilde{f}\). Note that random variable \(\alpha \nabla \tilde{f}\) has distribution \(p_\alpha(z) = \alpha^{-1} p(z/\alpha)\).

Now, suppose that \(\alpha\) is small enough that we can neglect the tails of \(p_\alpha(z)\) that lie outside the interval \([-1, 1]\). The probability of transitioning from \(x^0 = 0\) to \(x^1 = 1\) using stochastic rounding, denoted by \(T_\alpha(0, 1)\), is then
\[
T_\alpha(0, 1) \approx \int_0^1 z p_\alpha(z) dz = \frac{1}{\alpha} \int_0^\alpha z p(z/\alpha) dz
\]
\[
= \alpha \int_0^{1/\alpha} p(x) x dx \approx \alpha \int_0^{\infty} p(x) x dx,
\]
where the first approximation is because we neglected the unlikely case that \(\alpha \nabla \tilde{f} > 1\), and the second approximation appears because we added a small tail probability to the estimate. These approximations get more accurate for small \(\alpha\). We see that, assuming the tails of \(p\) are “light” enough, we have \(T_\alpha(0, 1) \sim \alpha \int_0^{\infty} p(x) x dx\) as \(\alpha \rightarrow 0\). Similarly, \(T_\alpha(0, -1) \sim \alpha \int_{-\infty}^{0} p(x) x dx\) as \(\alpha \rightarrow 0\).

What does this observation mean for the behavior of SR? First of all, the probability of leaving \(x^0\) on an iteration is
\[
T_\alpha(0, -1) + T_\alpha(0, 1)
\]
\[
\approx \alpha \left[ \int_0^{\infty} p(x) x dx + \int_{-\infty}^0 p(x) x dx \right], \tag{11}
\]
which vanishes for small \(\alpha\) (as also indicated by Lemma 1). This means the algorithm slows down as the learning rate drops off, which is not surprising. However, the conditional probability of ending up at \(x^1 = 1\) given that the algorithm did leave \(x^0\) is
\[
T_\alpha(0, 1|x^0 \neq x^0) \approx \frac{T_\alpha(0, 1)}{T_\alpha(0, -1) + T_\alpha(0, 1)}
\]
\[
= \frac{\int_0^{\infty} p(x) x dx}{\int_{-\infty}^0 p(x) x dx + \int_0^{\infty} p(x) x dx},
\]
which does not depend on \(\alpha\). In other words, provided \(\alpha\) is small, SR, on average, makes the same decisions/transitions with learning rate \(\alpha\) as it does with learning rate \(\alpha/10\); it just takes 10 times longer to make those
Theorem 2. Let \( p_{x,k} \) denote the probability distribution for the \( k \)th entry in \( \nabla \hat{f}(x) \), the stochastic gradient estimate at \( x \). Assume there is a constant \( C_1 \) such that for all \( x, k \), and \( v \) we have \( \int_{-\infty}^{\infty} p_{x,k}(z) \, dz \leq \frac{C_1}{\Delta} \), and some \( C_2 \) such that both \( \int_{-\infty}^{\infty} p_{x,k}(z) \, dz > 0 \) and \( \int_{-\infty}^{\infty} C_2 p_{x,k}(z) \, dz > 0 \). Define the matrix

\[
\tilde{U}(x,y) = \begin{cases} 
\int_{-\infty}^{\infty} p_{x,k}(z) \frac{z}{\Delta} \, dz, & \text{if } x - k = y - k, \ y_k = x_k + \Delta, \\
\int_{-\infty}^{\infty} C_2 p_{x,k}(z) \frac{z}{\Delta} \, dz, & \text{if } x - k = y - k, \ y_k = x_k - \Delta, \\
0, & \text{otherwise},
\end{cases}
\]

where \( x - k = y - k \) denotes that \( x \) and \( y \) are equal except at coordinate \( k \). Define the associated Markov chain transition matrix

\[
\tilde{T}_{\alpha_0} = I - \alpha_0 \cdot \text{diag}(1^T \tilde{U}) + \alpha_0 \tilde{U}, \tag{12}
\]

where \( \alpha_0 \) is the largest constant that makes \( \tilde{T}_{\alpha_0} \) non-negative. Suppose \( \tilde{T}_\alpha \) has a stationary distribution, denoted \( \tilde{\pi} \). Then, for sufficiently small \( \alpha \), \( T_\alpha \) has a stationary distribution \( \pi_\alpha \), and

\[
\lim_{\alpha \to 0} \pi_\alpha = \tilde{\pi}.
\]

Furthermore, this limiting distribution satisfies \( \tilde{\pi}(x) > 0 \) for any state \( x \), and is thus not concentrated on local minimizers of \( f \).

Proof. Let the matrix \( U_\alpha \) be a partial transition matrix defined by \( U_\alpha(x,x) = 0 \), and \( U_\alpha(x,y) = T_\alpha(x,y) \) for \( x \neq y \). From \( U_\alpha \), we can get back the full transition matrix \( T_\alpha \) using the formula

\[
T_\alpha = I - \text{diag}(1^T U_\alpha) + U_\alpha.
\]

Note that this formula is essentially “filling in” the diagonal entries of \( T_\alpha \) so that every column sums to 1, thus making \( T_\alpha \) a valid stochastic matrix.

Let’s bound the entries in \( U_\alpha \). Suppose that we begin an iteration of the stochastic rounding algorithm at some point
Consider an adjacent point \( y \) that differs from \( x \) at only 1 coordinate, \( k \), with \( y_k = x_k + \Delta \). Then we have

\[
U_\alpha(x, y) = \frac{1}{\alpha} \int_0^\Delta p_{x,k}(x/\alpha) \frac{x}{\Delta} \, dx + \frac{1}{\alpha} \int_0^{2\Delta} p_{x,k}(x/\alpha) \frac{2\Delta - x}{\Delta} \, dx
\]

\[
= \frac{1}{\alpha} \int_0^{\Delta/\alpha} p_{x,k}(z) \frac{z}{\alpha} \, dz + \frac{1}{\alpha} \int_{\Delta/\alpha}^{2\Delta/\alpha} p_{x,k}(z) \frac{2\Delta - az}{\Delta} \, dz
\]

\[
\leq \alpha \int_0^{\Delta/\alpha} p_{x,k}(z) \frac{z}{\Delta} \, dz + \int_{\Delta/\alpha}^{\infty} p_{x,k}(z) \, dz
\]

\[
= \alpha \int_0^\infty p_{x,k}(z) \frac{z}{\Delta} \, dz + O(\alpha^2). \tag{13}
\]

Note we have used the decay assumption:

\[
\int_\nu^\infty p_{x,k}(z) \leq \frac{C}{y^2}.
\]

Likewise, if \( y_k = x_k - \Delta \), then the transition probability is

\[
U_\alpha(x, y) = \alpha \int_{-\infty}^0 p_{x,k}(z) \frac{z}{\Delta} \, dz + O(\alpha^2), \tag{14}
\]

and if \( y_k = x_k \pm m\Delta \) for an integer \( m > 1 \),

\[
U_\alpha(x, y) = O(\alpha^2). \tag{15}
\]

We can approximate the behavior of \( U_\alpha \) using the matrix \( \hat{U}(x, y) \), and we define the associated markov chain transition matrix \( \hat{T}_\alpha \), as in (12).

Let \( \alpha_0 \) be the largest scalar such that the stochastic linear operator \( \hat{T}_{\alpha_0} \) has non-negative entries. For \( \alpha < \alpha_0 \), \( \hat{T}_\alpha \) has non-negative entries and column sums equal to 1; it thus defines the transition operator of a markov chain. Let \( \hat{\pi} \) denote the stationary distribution of the markov chain with transition matrix \( \hat{T}_{\alpha_0} \).

We now claim that \( \hat{\pi} \) is also the stationary distribution of \( \hat{T}_\alpha \) for all \( \alpha < \alpha_0 \). We verify this by noting that

\[
\hat{T}_\alpha = (I - \alpha \cdot \text{diag}(1^T \hat{U})) \hat{\pi} + \alpha \hat{U} \hat{\pi}
\]

\[
= (1 - \frac{\alpha}{\alpha_0}) I + \frac{\alpha}{\alpha_0} [I - \alpha_0 \cdot \text{diag}(1^T \hat{U}) + \alpha_0 \hat{U}]
\]

\[
= (1 - \frac{\alpha}{\alpha_0}) I + \frac{\alpha}{\alpha_0} \hat{T}_{\alpha_0}
\]

and so \( \hat{T}_\alpha \hat{\pi} = (1 - \frac{\alpha}{\alpha_0}) \hat{\pi} + \frac{\alpha}{\alpha_0} \hat{\pi} = \hat{\pi} \).

Recall that \( \hat{T}_\alpha \) is the transition matrix for the Markov chain generated by the stochastic rounding algorithm with learning rate \( \alpha \). We wish to show that this markov chain is well approximated by \( \hat{T}_\alpha \). Note that

\[
T_\alpha(x, y) = \prod_{k, x_k \neq y_k} T_{\alpha}(x, x + (y_k - x_k) \Delta e_k) \leq O(\alpha^2)
\]

when \( x, y \) differ at more than 1 coordinate. In other words, transitions between multiple coordinates simultaneously become vanishingly unlikely for small \( \alpha \). When \( x \) and \( y \) differ by exactly 1 coordinate, we know from (13) that

\[
T_\alpha(x, y) = \alpha U(x, y) + O(\alpha^2).
\]

These observations show that the off-diagonal elements of \( T_\alpha \) are well approximated (up to uniform \( O(\alpha^2) \) error) by the corresponding elements in \( \alpha U \). Since the columns of \( T_\alpha \) sum to one, the diagonal elements are well approximated as well, and we have

\[
T_\alpha = (I - \alpha \cdot \text{diag}(1^T U)) + \alpha U + O(\alpha^2) = \hat{T}_\alpha + O(\alpha^2).
\]

To be precise, the notation above means that

\[
|T_\alpha(x, y) - \hat{T}_\alpha(x, y)| < C\alpha^2, \tag{17}
\]

for some \( C \) that is uniform over \( (x, y) \).

We are now ready to show that the stationary distribution of \( T_\alpha \) exists and approaches \( \hat{\pi} \). Re-arranging (16) gives us

\[
\alpha_0 \hat{T}_\alpha + (\alpha - \alpha_0) I = \alpha \hat{T}_{\alpha_0}.
\]

Combining this with (17), we get

\[
\|\alpha_0 T_\alpha + (\alpha - \alpha_0) I - \alpha \hat{T}_{\alpha_0}\|_\infty < O(\alpha^2),
\]

and so

\[
\|\frac{\alpha_0}{\alpha} T_\alpha + (1 - \frac{\alpha_0}{\alpha}) I - \hat{T}_{\alpha_0}\|_\infty < O(\alpha). \tag{18}
\]

From (18), we see that the matrix \( \frac{\alpha_0}{\alpha} T_\alpha + (1 - \frac{\alpha_0}{\alpha}) I \) approaches \( \hat{T}_{\alpha_0} \). Note that \( \hat{\pi} \) is the Perron-Frobenius eigenvalue of \( \hat{T}_{\alpha_0} \), and thus has multiplicity 1. Multiplicity 1 eigenvalues/vectors of a matrix vary continuously with small perturbations to that matrix (Theorem 8, p130 of (Lax, 2007)). It follows that, for small \( \alpha_0 \), \( \frac{\alpha_0}{\alpha} T_\alpha + (1 - \frac{\alpha_0}{\alpha}) I \) has a stationary distribution, and this distribution approaches \( \hat{\pi} \). The leading eigenvector of \( \frac{\alpha_0}{\alpha} T_\alpha + (1 - \frac{\alpha_0}{\alpha}) I \) is the same as the leading eigenvector of \( T_\alpha \), and it follows that \( T_\alpha \) has a stationary distribution that approaches \( \hat{\pi} \).

Finally, note that we have assumed \( \int_0^{C_2} p_{x,k}(z) \, dz > 0 \) and \( \int_{-C_2}^0 p_{x,k}(z) \, dz > 0 \). Under this assumption, for \( \alpha < \frac{1}{C_2} \), \( \hat{T}_{\alpha_0}(x, y) > 0 \) whenever \( x, y \) are neighbors differ at a single coordinate. It follows that every state in the Markov chain \( \hat{T}_{\alpha_0} \) is accessible from every other state by traversing a path of non-zero transition probabilities, and so \( \hat{\pi}(x) > 0 \) for every state \( x \).

While the long term stationary behavior of SR is relatively insensitive to \( \alpha \), the convergence speed of the algorithm is not. To measure this, we consider the mixing time of the Markov chain. Let \( \pi_\alpha \) denote the stationary distribution of a Markov chain. We say that the \( \epsilon \)-mixing time of the chain is \( M_\epsilon \) if \( M_\epsilon \) is the smallest integer such that (Levin et al., 2009)

\[
|\mathbb{P}(X^{M_\epsilon} \in A|x^0) - \pi(A)| \leq \epsilon, \forall x^0, \forall \text{ subsets } A \subseteq X. \tag{19}
\]
Towards a Deeper Understanding of Training Quantized Neural Networks

We show below that the mixing time of the Markov chain gets large for small $\alpha$, which means exploration slows down, even though no exploitation gain is being realized.

**Theorem 3.** Let $p_{x,k}$ satisfy the assumptions of Theorem 2. Choose some $\epsilon$ sufficiently small that there exists a proper subset of states $A \subset X$ with stationary probability $\pi_A(\alpha)$ greater than $\epsilon$. Let $M(\alpha)$ denote the $\epsilon$-mixing time of the chain with learning rate $\alpha$. Then,

$$\lim_{\alpha \to 0} M(\alpha) = \infty.$$

**Proof.** Given some distribution $\pi$ over the states of the markov chain, and some set $A$ of states, let $[\pi]_A = \sum_{a \in A} \pi(a)$ denote the measure of $A$ with respect to $\pi$.

Suppose for contradiction that the mixing time of the chain remains bounded as $\alpha$ vanishes. Then we can find an integer $M$, that upper bounds the $\epsilon$-mixing time for all $\alpha$. By the assumption of the theorem, we can select some set of states $A$ with $[\pi]_A > \epsilon$, and some starting state $y \notin A$. Let $e$ be a distribution (a vector in the finite-state case) with $e_y = 1, e_k = 0$ for $k \neq y$. Note that $[e]_A = 0$ because $y \notin A$. Then

$$|[e]_A - [\pi]_A| > \epsilon.$$

Note that, as $\alpha \to 0$, we have $\|T_\alpha - \tilde{T}_\alpha\| \to 0$ and thus $\|T^{M}_\alpha - \tilde{T}^{M}_\alpha\| \to 0$. We also see from the definition of $\tilde{T}_\alpha$ in (12), $\lim_{\alpha \to 0} \tilde{T}_\alpha = I$. It follows that

$$\lim_{\alpha \to 0} \|T^{M}_\alpha e - [\pi]_A\| = \|e - [\pi]_A\| > \epsilon,$$

and so for some $\alpha$ the inequality (19) is violated. This is a contradiction because it was assumed $M$ is an upper bound on the mixing time. \qed

6. Experiments

To explore the implications of the theory above, we train two types of networks, VGG-like CNNs (Simonyan & Zisserman, 2015) and Residual networks (He et al., 2016), on CIFAR-10 and CIFAR-100. VGG-BC is a high-capacity network used for the original BC model (Courbariaux et al., 2015). We use the same architecture as in (Courbariaux et al., 2015) except using softmax and cross-entropy loss instead of SVM and squared hinge loss, respectively. ResNets-56 has 55 convolutional layers and one linear layer, and contains three stages of residual blocks where each stage has the same number of residual blocks. We also create a wide ResNet-56 (WRN-56-2) that doubles the number of filters in each residual block as in (Zagoruyko & Komodakis, 2016). We use ADAM (Kingma & Ba, 2015) as the baseline optimizer, and train with the three quantized algorithms mentioned in Section 3, i.e., R-ADAM, SR-ADAM and BC-ADAM. All methods start with the same binary weight initialization of random $\pm 1$s. The image pre-processing and data augmentation procedures are the same as (He et al., 2016). Similar to (Rastegari et al., 2016), here we only quantize the weights in the convolutional layers, but not linear layers, during training. We set the initial learning rate to 0.01 and decrease the learning rate by a factor of 10 at epochs 82 and 122 for CIFAR-10 and CIFAR-100 (He et al., 2016). The default minibatch size is 128. Following (Courbariaux et al., 2015), we do not use weight decay during training. We implement all quantization algorithms in Torch7 (Collobert et al., 2011) and train the quantized models with NVIDIA GPUs.

**Results** The overall results are summarized in Table 1. The binary model trained by BC-ADAM has comparable performance to the full-precision model trained by ADAM. SR-ADAM outperforms R-ADAM, which verifies the effectiveness of Stochastic Rounding. There is a performance gap between SR-ADAM and BC-ADAM across all models and datasets. This is consistent with our theoretical results in Sections 4 and 5, and keeping track of the real-valued weights to quantize the binary weights in BC-ADAM seems to really help empirically.

7. Conclusion

The training of quantized neural nets is essential for deploying machine learning models on embedded and ubiquitous devices. In this work, we provided a theoretical analysis to better understand the BinaryConnect (BC) and Stochastic Rounding (SR) methods for training binary nets. We showed that BC algorithms have the desirable property of being able to concentrate on minimizers. For general non-convex problems, however, we proved that SR differs from conventional stochastic methods in that it is unable to exploit greedy local search. Experiments confirm these findings, and show that the mathematical properties of SR are indeed observable (and very important) in practice.

**Acknowledgments** T. Goldstein was supported in part by the US National Science Foundation (NSF) under grant CCF-1535902, by the US Office of Naval Research under grant N00014-17-1-2078, and by the Sloan Foundation. C. Studer was supported in part by Xilinx, Inc. and by the US NSF under grants ECCS-1408006, CCF-1535897, and CA-REER CCF-1652065. H. Samet was supported in part by the US NSF under grant IIS-13-20791.
References


Hwang, Kyuyeon and Sung, Wonyong. Fixed-point feedforward deep neural network design using weights+ 1, 0, and -1. In IEEE Workshop on Signal Processing Systems (SiPS), 2014.


Lin, Zhouran, Courbariaux, Matthieu, Memisevic, Roland, and Bengio, Yoshua. Neural networks with few multiplications. ICLR, 2016b.


Zhu, Chenzhuo, Han, Song, Mao, Huizi, and Dally, William J. Trained ternary quantization. ICLR, 2017.