## Tarcil: Reconciling Scheduling Speed and Quality in Large Shared Clusters

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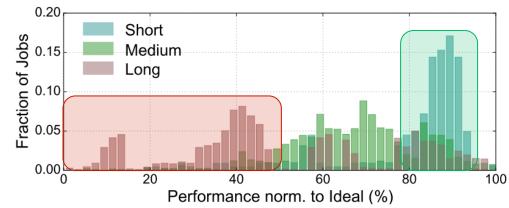
*SOCC – August 27<sup>th</sup> 2015* 

## **Executive Summary**

- Goals of cluster scheduling
  - High decision quality
  - High scheduling speed
- High performance
   High cluster utilization
- Problem: Disparity in scheduling designs
  - Centralized schedulers  $\rightarrow$  High quality, low speed
  - Sampling-based schedulers  $\rightarrow$  High speed, low quality
- Tarcil: Key scheduling techniques to bridge the gap
  - Account for resource preferences  $\rightarrow$  High decision quality
  - Analytical framework for sampling  $\rightarrow$  Predictable performance
  - □ Admission control  $\rightarrow$  High quality & speed
  - Distributed design  $\rightarrow$  High scheduling speed

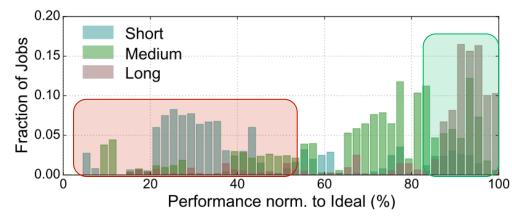
#### Motivation

Optimize scheduling speed (sampling-based, distributed)



#### Good: Short jobs Bad: Long jobs

Optimize scheduling quality (centralized, greedy)



#### Good: Long jobs Bad: Short jobs

Short: 100msec, Medium: 1-10sec, Long: 10sec-10min

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## Key Scheduling Techniques at Scale

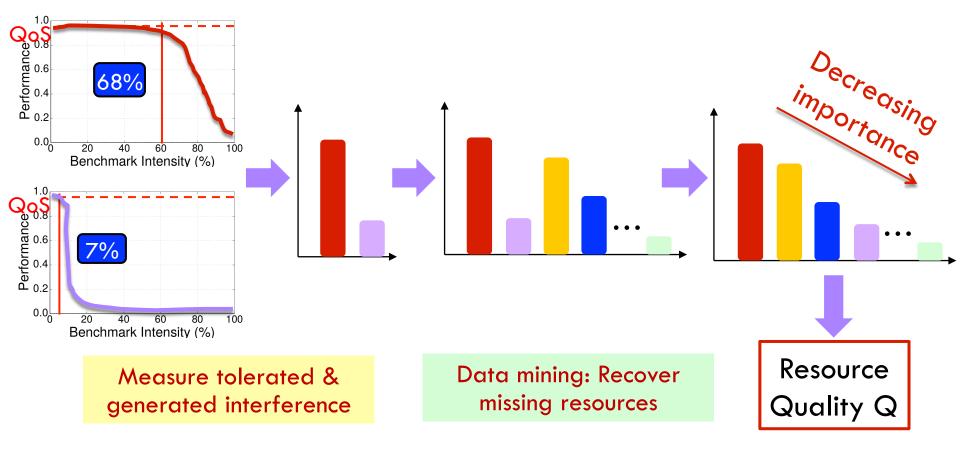
## 1. Determine Resource Preferences

- Scheduling quality depends on: interference, heterogeneity, scale up/out, ...
  - **\square** Exhaustive exploration  $\rightarrow$  infeasible
  - Practical data mining framework<sup>1</sup>
  - Measure impact of a couple of allocations 

     A stimate for large space

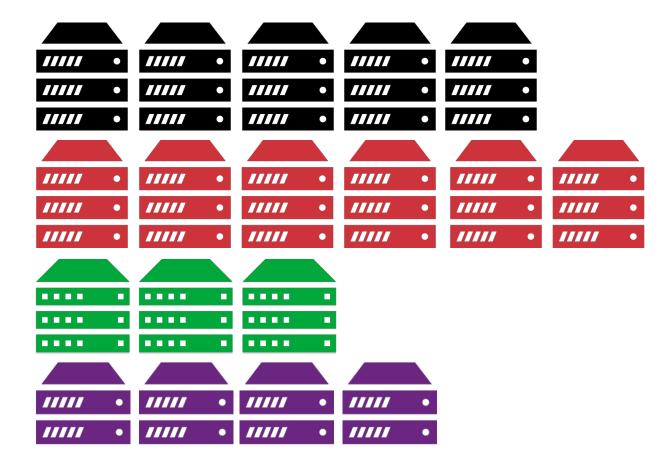
## **Example: Quantifying Interference**

Interference: set of microbenchmarks of tunable intensity (iBench)



<sup>1</sup>C. Delimitrou and C. Kozyrakis. Quasar: Resource-Efficient and QoS-Aware Cluster Management. In ASPLOS 2014.

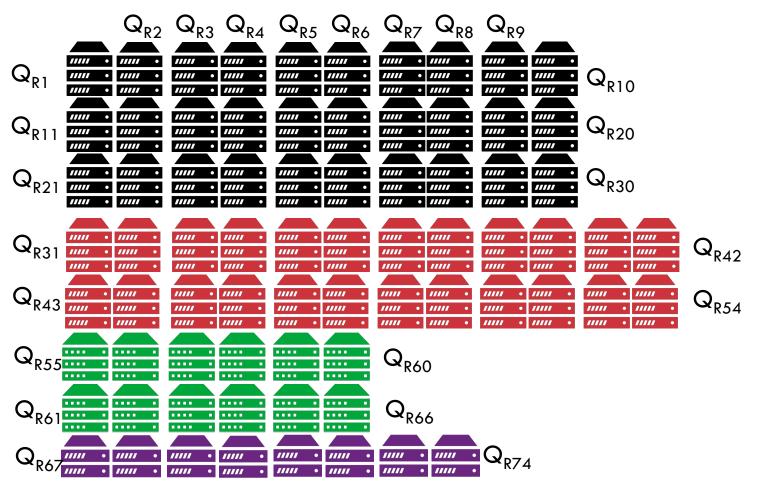
Sample w.r.t. required resource quality



 $\hfill \label{eq:constraint}$  Fine-grain allocations: partition servers in Resource Units (RU)  $\rightarrow$  minimum allocation unit

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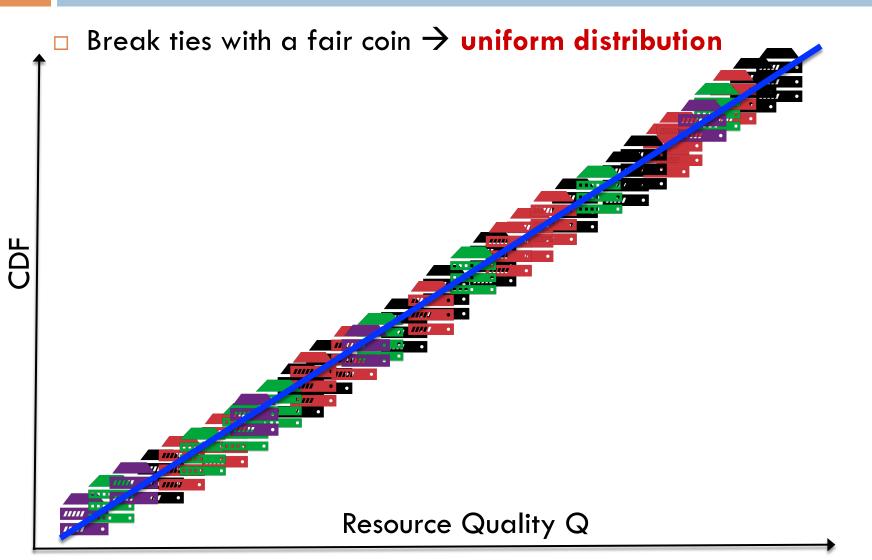
Match a new job with required quality Q to appropriate RUs

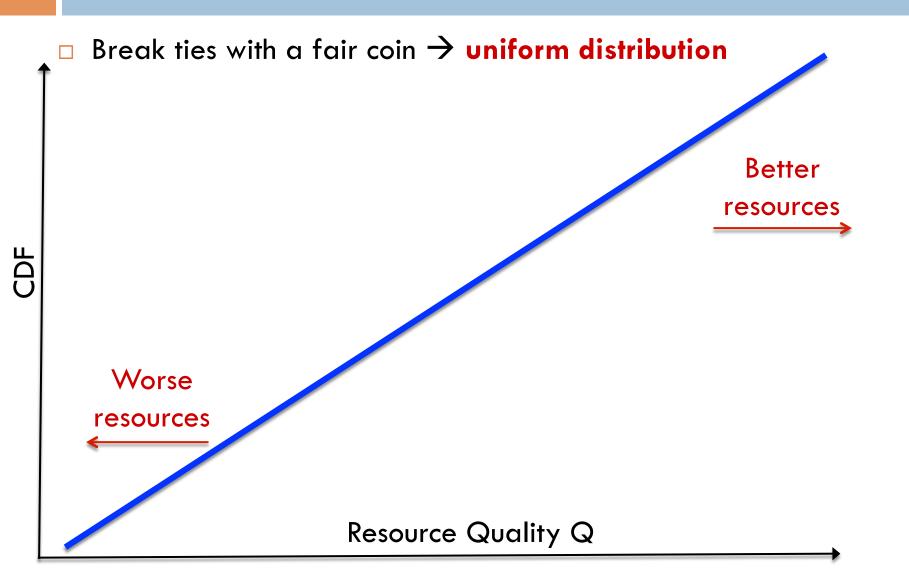


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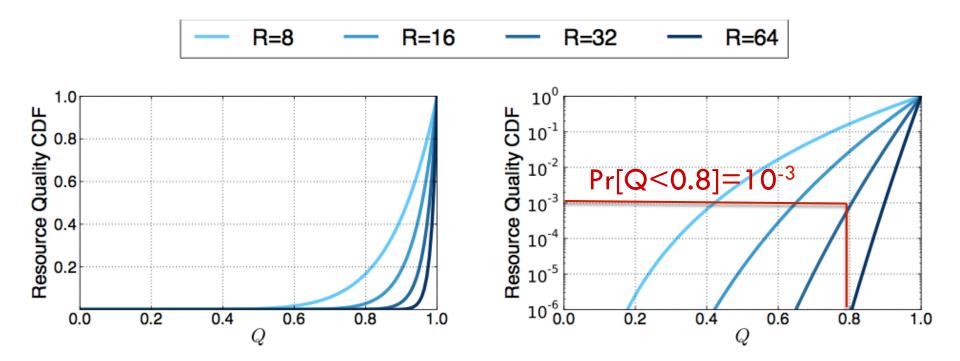
#### Rank resources by quality

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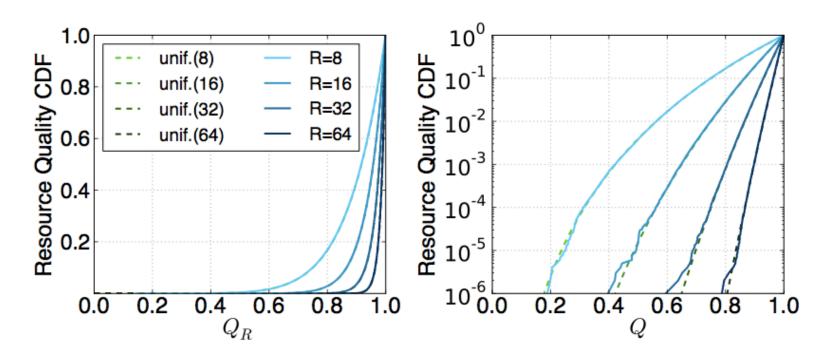




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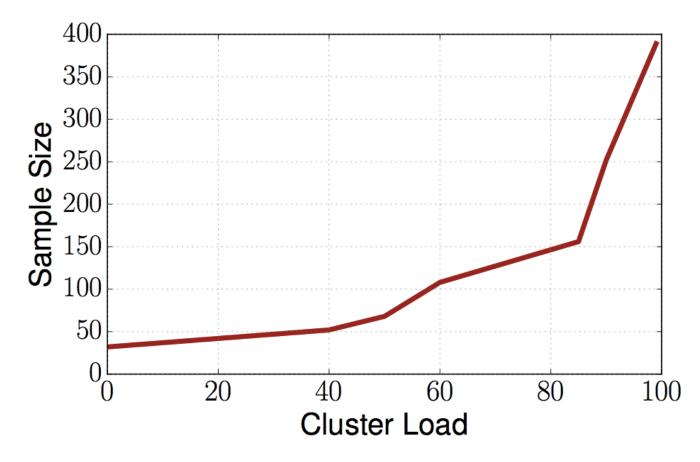
# Validation



- 100 server EC2 cluster
- Short Spark tasks
- Deviation between analytical and empirical is minimal

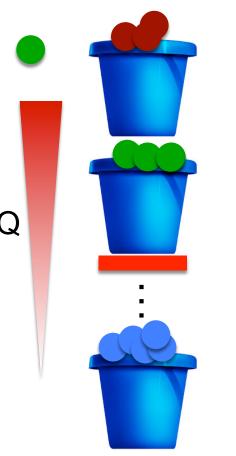
# Sampling at High Load

- Performance degrades (with small sample size)
- Or sample size needs to increase



## 3. Admission Control

- Queue jobs based on required resource quality
- $\square$  Resource quality vs. waiting time  $\rightarrow$  set max waiting time limit



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## **Tarcil Implementation**

- $\Box$  4,000 loc in C/C++ and Python
- Supports apps in various frameworks (Hadoop, Spark, key-value stores)
- Distributed design: Concurrent scheduling agents (sim. Omega<sup>2</sup>)
  - Each agent has local copy of state, one resilient master copy
  - Lock-free optimistic concurrency for conflict resolution (rare) → Abort and retry
  - 30:1 worker to scheduling agent ratio

<sup>2</sup>M. Schwarzkopf, A. Konwinski, et al. Omega: flexible, scalable schedulers for large compute clusters. In EuroSys 2013.

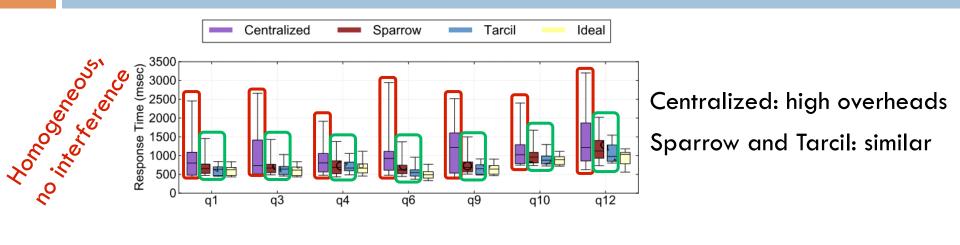
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## **Evaluation Methodology**

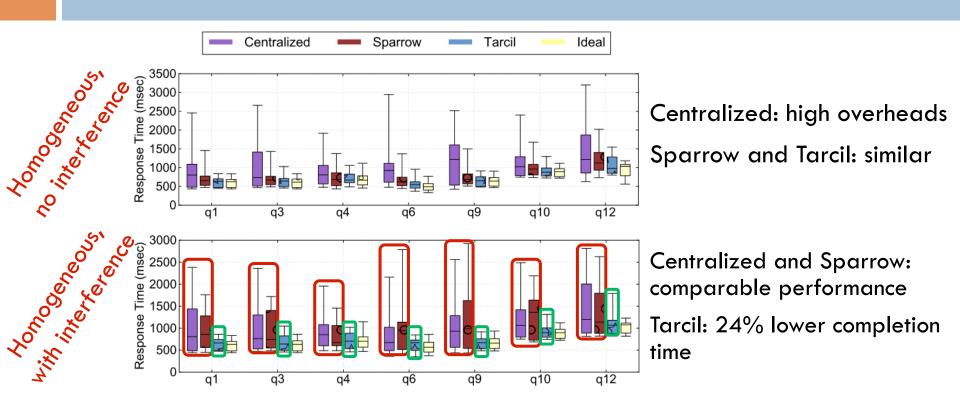
#### 1. TPC-H Workload

- ~40k queries of different types
- Compare with a centralized scheduler (Quasar) and a distributed scheduler based on random sampling (Sparrow)
- 110-server EC2 cluster (100 workers, 10 scheduling agents)
  - Homogeneous cluster, no interference
  - Homogeneous cluster, with interference
  - Heterogeneous cluster, with interference
- Metrics:
  - Task performance
  - Performance predictability
  - Scheduling latency

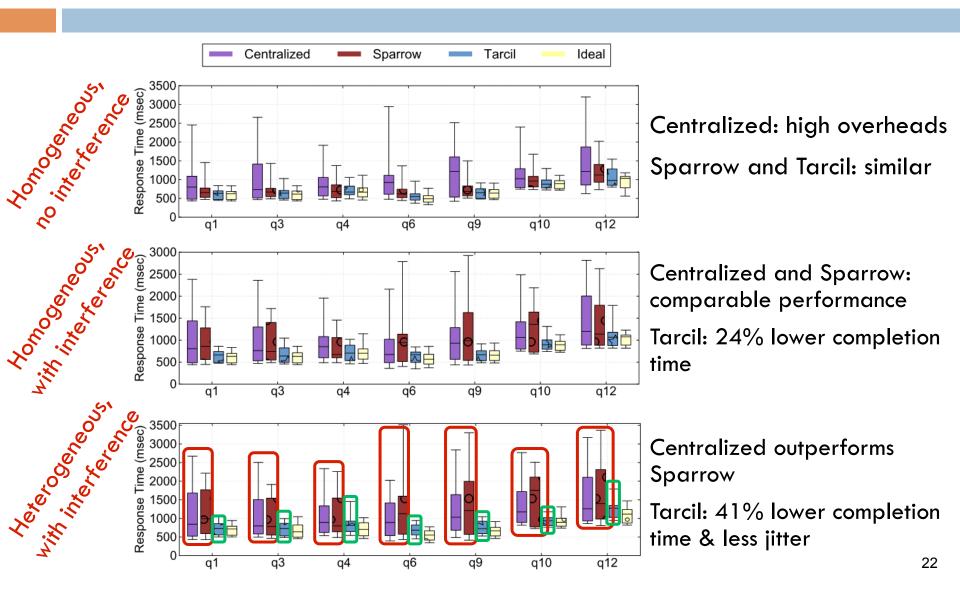
## Evaluation



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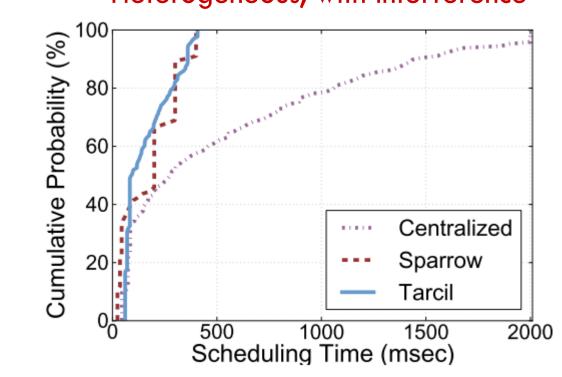


## Evaluation



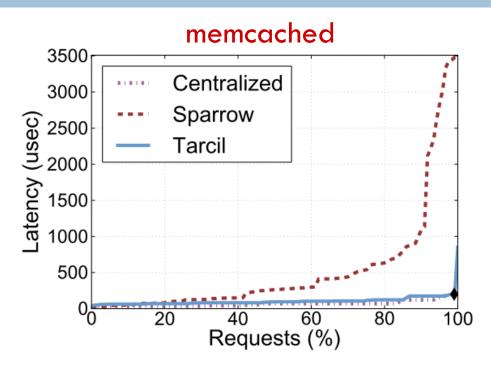
## Scheduling Overheads

Heterogeneous, with interference



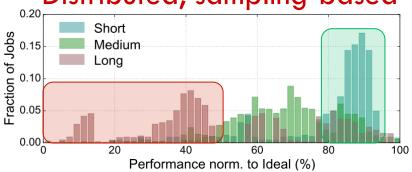
- Centralized: Two orders of magnitude slower than the distributed, sampling-based schedulers
- □ Sparrow and Tarcil: Comparable scheduling overheads

## **Resident Load**

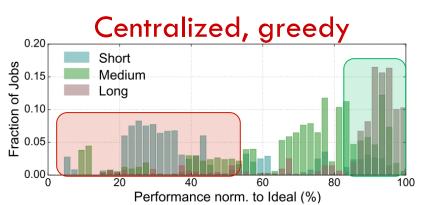


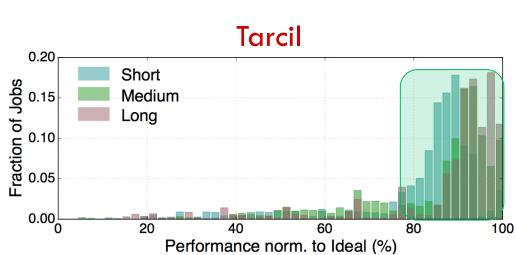
- □ Tarcil and Centralized account for cross-job interference → preserve memcached's QoS
- Sparrow causes QoS violations for memcached

## **Motivation Revisited**









Short: 100msec Medium: 1-10sec Long:10sec-10min

## More details in the paper...

#### Sensitivity on parameters such as:

- Cluster load
- Number of scheduling agents
- Sample size
- Task duration, etc.
- Job priorities
- Large allocations
- Generic application scenario (batch and latency-critical) on 200 EC2 servers

## Conclusions

Tarcil: Reconciles high quality and high speed scheduling

- Account for resource preferences
- Analytical sampling framework to improve predictability
- Admission control to maintain high scheduling quality at high load
- Distributed design to improve scheduling speed

Results:

- 41% better performance than random sampling-based schedulers
- 100x better scheduling latency than centralized schedulers
- Predictable allocation quality & performance

## Questions?

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#### Questions??

#### Thank you